

Internet Engineering Task Force (IETF)
Request for Comments: 8029
Obsoletes: 4379, 6424, 6829, 7537
Updates: 1122
Category: Standards Track
ISSN: 2070-1721

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March 2017

Detecting Multiprotocol Label Switched (MPLS) Data-Plane Failures

Abstract

This document describes a simple and efficient mechanism to detect data-plane failures in Multiprotocol Label Switching (MPLS) Label Switched Paths (LSPs). It defines a probe message called an "MPLS echo request" and a response message called an "MPLS echo reply" for returning the result of the probe. The MPLS echo request is intended to contain sufficient information to check correct operation of the data plane and to verify the data plane against the control plane, thereby localizing faults.

This document obsoletes RFCs 4379, 6424, 6829, and 7537, and updates RFC 1122.

Status of This Memo

This is an Internet Standards Track document.

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1. Introduction

This document describes a simple and efficient mechanism to detect data-plane failures in MPLS Label Switched Paths (LSPs). It defines a probe message called an "MPLS echo request" and a response message called an "MPLS echo reply" for returning the result of the probe. The MPLS echo request is intended to contain sufficient information to check correct operation of the data plane, as well as a mechanism to verify the data plane against the control plane, thereby localizing faults.

An important consideration in this design is that MPLS echo requests follow the same data path that normal MPLS packets would traverse. MPLS echo requests are meant primarily to validate the data plane and secondarily to verify the data plane against the control plane. Mechanisms to check the control plane are valuable but are not covered in this document.

This document makes special use of the address range 127/8. This is an exception to the behavior defined in RFC 1122 [RFC1122], and this specification updates that RFC. The motivation for this change and the details of this exceptional use are discussed in Section 2.1 below.

This document obsoletes RFC 4379 [RFC4379], RFC 6424 [RFC6424], RFC 6829 [RFC6829], and RFC 7537 [RFC7537].

1.1. Conventions

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119 [RFC2119].

The term "Must Be Zero" (MBZ) is used in object descriptions for reserved fields. These fields MUST be set to zero when sent and ignored on receipt.

Terminology pertaining to L2 and L3 Virtual Private Networks (VPNs) is defined in [RFC4026].

Since this document refers to the MPLS Time to Live (TTL) far more frequently than the IP TTL, the authors have chosen the convention of using the unqualified "TTL" to mean "MPLS TTL" and using "IP TTL" for the TTL value in the IP header.

1.2. Structure of This Document

The body of this memo contains four main parts: motivation, MPLS echo request/reply packet format, LSP ping operation, and a reliable return path. It is suggested that first-time readers skip the actual packet formats and read the "Theory of Operation" (Section 4) first; the document is structured the way it is to avoid forward references.

1.3. Scope of This Specification

The primary goal of this document is to provide a clean and updated LSP ping specification.

[RFC4379] defines the basic mechanism for MPLS LSP validation that can be used for fault detection and isolation. The scope of this document also includes various updates to MPLS LSP ping, including:

- o Update all references and citations.
 - * Obsoleted RFCs 2434, 2030, and 3036 are respectively replaced with RFCs 5226, 5905, and 5036.
 - * Additionally, some informative references were published as RFCs: RFCs 4761, 5085, 5885, and 8077.
- o Incorporate all outstanding RFC errata.
 - * See [Err108], [Err742], [Err1418], [Err1714], [Err1786], [Err2978], [Err3399].
- o Replace EXP with Traffic Class (TC), based on the update from RFC 5462.
- o Incorporate the updates from RFC 6829, by adding the pseudowire (PW) Forwarding Equivalence Classes (FECs) advertised over IPv6 and obsoleting RFC 6829.
- o Incorporate the updates from RFC 7506, by adding the IPv6 Router Alert Option (RAO) for MPLS Operations, Administration, and Maintenance (OAM).
- o Incorporate newly defined bits on the Global Flags field from RFCs 6425 and 6426.
- o Update the IPv4 addresses used in examples to utilize the documentation prefix. Add examples with IPv6 addresses.

- o Incorporate the updates from RFC 6424, by deprecating the Downstream Mapping TLV (DSMAP) and adding the Downstream Detailed Mapping TLV (DDMAP); updating two new Return Codes; adding the motivations of tunneled or stitched LSPs; updating the procedures, IANA considerations, and security considerations; and obsoleting RFC 6424.
- o Incorporate the updates from RFC 7537, by updating the IANA Considerations section and obsoleting RFC 7537.
- o Finally, obsolete RFC 4379.

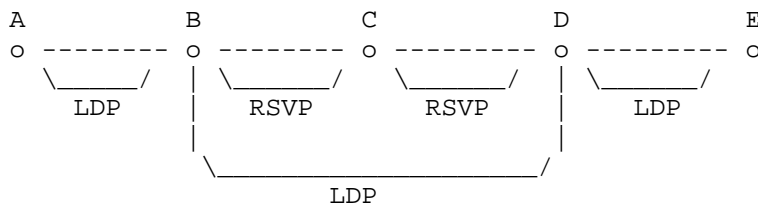
2. Motivation

When an LSP fails to deliver user traffic, the failure cannot always be detected by the MPLS control plane. There is a need to provide a tool that would enable users to detect such traffic "black holes" or misrouting within a reasonable period of time and a mechanism to isolate faults.

In this document, we describe a mechanism that accomplishes these goals. This mechanism is modeled after the ping/traceroute paradigm: ping (ICMP echo request [RFC0792]) is used for connectivity checks, and traceroute is used for hop-by-hop fault localization as well as path tracing. This document specifies a "ping" mode and a "traceroute" mode for testing MPLS LSPs.

The basic idea is to verify that packets that belong to a particular FEC actually end their MPLS path on a Label Switching Router (LSR) that is an egress for that FEC. This document proposes that this test be carried out by sending a packet (called an "MPLS echo request") along the same data path as other packets belonging to this FEC. An MPLS echo request also carries information about the FEC whose MPLS path is being verified. This echo request is forwarded just like any other packet belonging to that FEC. In "ping" mode (basic connectivity check), the packet should reach the end of the path, at which point it is sent to the control plane of the egress LSR, which then verifies whether it is indeed an egress for the FEC. In "traceroute" mode (fault isolation), the packet is sent to the control plane of each transit LSR, which performs various checks to confirm that it is indeed a transit LSR for this path; this LSR also returns further information that helps check the control plane against the data plane, i.e., that forwarding matches what the routing protocols determined as the path.

An LSP traceroute may cross a tunneled or stitched LSP en route to the destination. While performing end-to-end LSP validation in such scenarios, the FEC information included in the packet by the Initiator may be different from the one assigned by the transit node in a different segment of a stitched LSP or tunnel. Let us consider a simple case.



When an LSP traceroute is initiated from Router A to Router E, the FEC information included in the packet will be LDP while Router C along the path is a pure RSVP node and does not run LDP. Consequently, node C will be unable to perform FEC validation. The MPLS echo request should contain sufficient information to allow any transit node within a stitched or tunneled LSP to perform FEC validations to detect any misrouted echo requests.

One way these tools can be used is to periodically ping a FEC to ensure connectivity. If the ping fails, one can then initiate a traceroute to determine where the fault lies. One can also periodically traceroute FECs to verify that forwarding matches the control plane; however, this places a greater burden on transit LSRs and thus should be used with caution.

2.1. Use of Address Range 127/8

As described above, LSP ping is intended as a diagnostic tool. It is intended to enable providers of an MPLS-based service to isolate network faults. In particular, LSP ping needs to diagnose situations where the control and data planes are out of sync. It performs this by routing an MPLS echo request packet based solely on its label stack. That is, the IP destination address is never used in a forwarding decision. In fact, the sender of an MPLS echo request packet may not know, a priori, the address of the router at the end of the LSP.

Providers of MPLS-based services also need the ability to trace all of the possible paths that an LSP may take. Since most MPLS services are based on IP unicast forwarding, these paths are subject to Equal-Cost Multipath (ECMP) load sharing.

This leads to the following requirements:

1. Although the LSP in question may be broken in unknown ways, the likelihood of a diagnostic packet being delivered to a user of an MPLS service MUST be held to an absolute minimum.
2. If an LSP is broken in such a way that it prematurely terminates, the diagnostic packet MUST NOT be IP forwarded.
3. A means of varying the diagnostic packets such that they exercise all ECMP paths is thus REQUIRED.

Clearly, using general unicast addresses satisfies neither of the first two requirements. A number of other options for addresses were considered, including a portion of the private address space (as determined by the network operator) and the IPv4 link-local addresses. Use of the private address space was deemed ineffective since the leading MPLS-based service is an IPv4 VPN. VPNs often use private addresses.

The IPv4 link-local addresses are more attractive in that the scope over which they can be forwarded is limited. However, if one were to use an address from this range, it would still be possible for the first recipient of a diagnostic packet that "escaped" from a broken LSP to have that address assigned to the interface on which it arrived and thus could mistakenly receive such a packet. Older deployed routers may not (correctly) implement IPv4 link-local addresses and would forward a packet with an address from that range toward the default route.

The 127/8 range for IPv4 and that same range embedded in an IPv4-mapped IPv6 address for IPv6 was chosen for a number of reasons.

RFC 1122 allocates the 127/8 as the "Internal host loopback address" and states: "Addresses of this form MUST NOT appear outside a host." Thus, the default behavior of hosts is to discard such packets. This helps to ensure that if a diagnostic packet is misdirected to a host, it will be silently discarded.

RFC 1812 [RFC1812] states:

A router SHOULD NOT forward, except over a loopback interface, any packet that has a destination address on network 127. A router MAY have a switch that allows the network manager to disable these checks. If such a switch is provided, it MUST default to performing the checks.

This helps to ensure that diagnostic packets are never IP forwarded.

The 127/8 address range provides 16M addresses allowing wide flexibility in varying addresses to exercise ECMP paths. Finally, as an implementation optimization, the 127/8 range provides an easy means of identifying possible LSP packets.

2.2. Router Alert Option

This document requires the use of the RAO set in an IP header in order to have the transit node process the MPLS OAM payload.

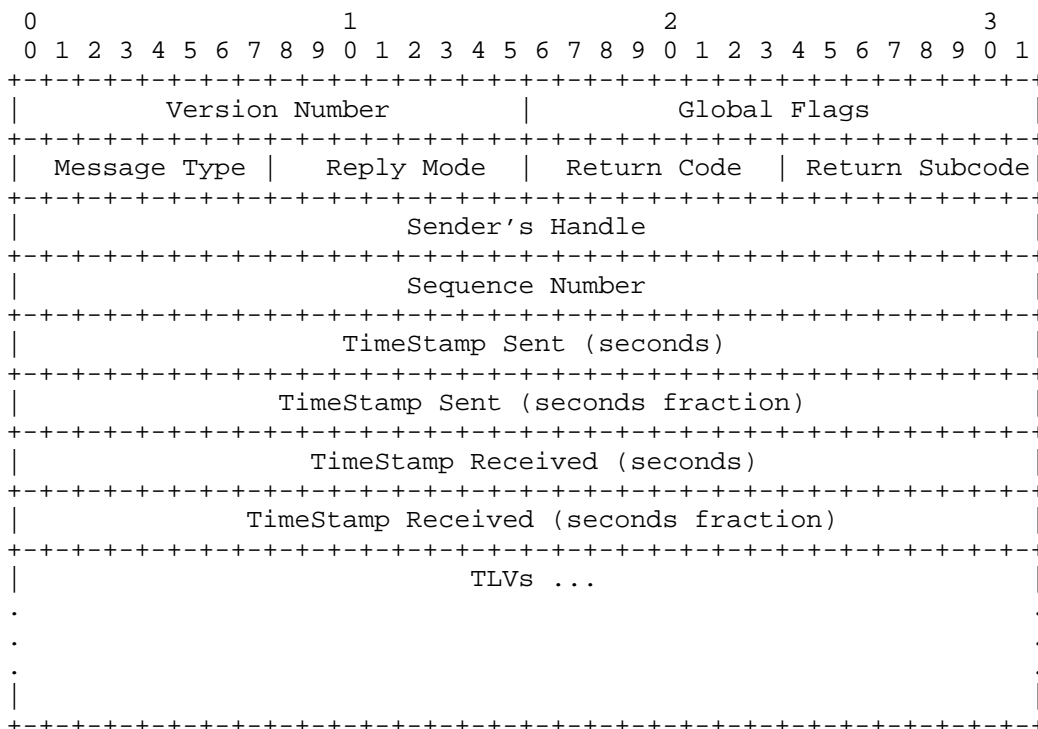
[RFC2113] defines a generic Option Value 0x0 for IPv4 RAO that alerts the transit router to examine the IPv4 packet. [RFC7506] defines MPLS OAM Option Value 69 for IPv6 RAO to alert transit routers to examine the IPv6 packet more closely for MPLS OAM purposes.

The use of the Router Alert IP Option in this document is as follows:

In case of an IPv4 header, the generic IPv4 RAO value 0x0 [RFC2113] SHOULD be used. In case of an IPv6 header, the IPv6 RAO value (69) for MPLS OAM [RFC7506] MUST be used.

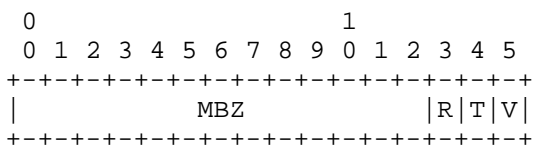
3. Packet Format

An MPLS echo request/reply is a (possibly labeled) IPv4 or IPv6 UDP packet; the contents of the UDP packet have the following format:



The Version Number is currently 1. (Note: the version number is to be incremented whenever a change is made that affects the ability of an implementation to correctly parse or process an MPLS echo request/reply. These changes include any syntactic or semantic changes made to any of the fixed fields, or to any Type-Length-Value (TLV) or sub-TLV assignment or format that is defined at a certain version number. The version number may not need to be changed if an optional TLV or sub-TLV is added.)

The Global Flags field is a bit vector with the following format:



At the time of writing, three flags are defined: the R, T, and V bits; the rest MUST be set to zero when sending and ignored on receipt.

The V (Validate FEC Stack) flag is set to 1 if the sender wants the receiver to perform FEC Stack validation; if V is 0, the choice is left to the receiver.

The T (Respond Only If TTL Expired) flag MUST be set only in the echo request packet by the sender. If the T flag is set to 1 in an incoming echo request, and the TTL of the incoming MPLS label is more than 1, then the receiving node MUST drop the incoming echo request and MUST NOT send any echo reply to the sender. This flag MUST NOT be set in the echo reply packet. If this flag is set in an echo reply packet, then it MUST be ignored. The T flag is defined in Section 3.4 of [RFC6425].

The R (Validate Reverse Path) flag is defined in [RFC6426]. When this flag is set in the echo request, the Responder SHOULD return reverse-path FEC information, as described in Section 3.4.2 of [RFC6426].

The Message Type is one of the following:

Value	Meaning
1	MPLS Echo Request
2	MPLS Echo Reply

The Reply Mode can take one of the following values:

Value	Meaning
1	Do not reply
2	Reply via an IPv4/IPv6 UDP packet
3	Reply via an IPv4/IPv6 UDP packet with Router Alert
4	Reply via application-level control channel

An MPLS echo request with 1 (Do not reply) in the Reply Mode field may be used for one-way connectivity tests; the receiving router may log gaps in the Sequence Numbers and/or maintain delay/jitter statistics. An MPLS echo request would normally have 2 (Reply via an IPv4/IPv6 UDP packet) in the Reply Mode field. If the normal IP return path is deemed unreliable, one may use 3 (Reply via an IPv4/IPv6 UDP packet with Router Alert). Note that this requires that all intermediate routers understand and know how to forward MPLS echo replies. The echo reply uses the same IP version number as the received echo request, i.e., an IPv4 encapsulated echo reply is sent in response to an IPv4 encapsulated echo request.

Some applications support an IP control channel. One such example is the associated control channel defined in Virtual Circuit Connectivity Verification (VCCV) [RFC5085][RFC5885]. Any application that supports an IP control channel between its control entities may set the Reply Mode to 4 (Reply via application-level control channel) to ensure that replies use that same channel. Further definition of this code point is application specific and thus beyond the scope of this document.

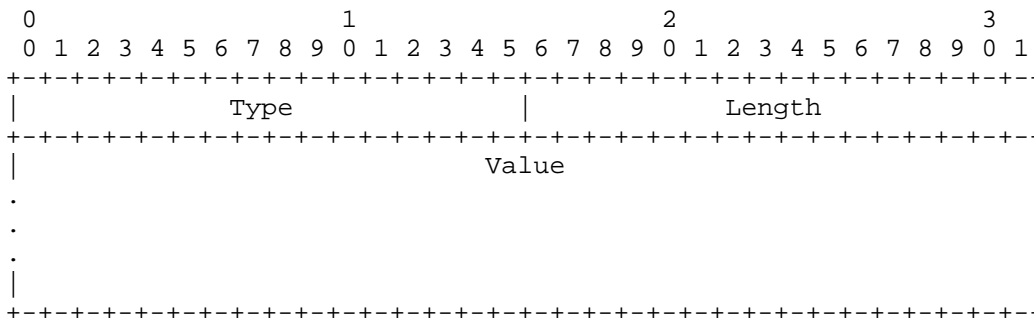
Return Codes and Subcodes are described in Section 3.1.

The Sender's Handle is filled in by the sender and returned unchanged by the receiver in the echo reply (if any). There are no semantics associated with this handle, although a sender may find this useful for matching up requests with replies.

The Sequence Number is assigned by the sender of the MPLS echo request and can be (for example) used to detect missed replies.

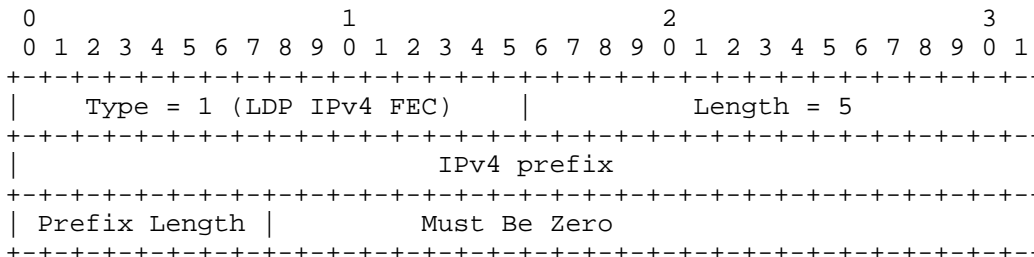
The TimeStamp Sent is the time of day (according to the sender's clock) in 64-bit NTP timestamp format [RFC5905] when the MPLS echo request is sent. The TimeStamp Received in an echo reply is the time of day (according to the receiver's clock) in 64-bit NTP timestamp format in which the corresponding echo request was received.

TLVs (Type-Length-Value tuples) have the following format:

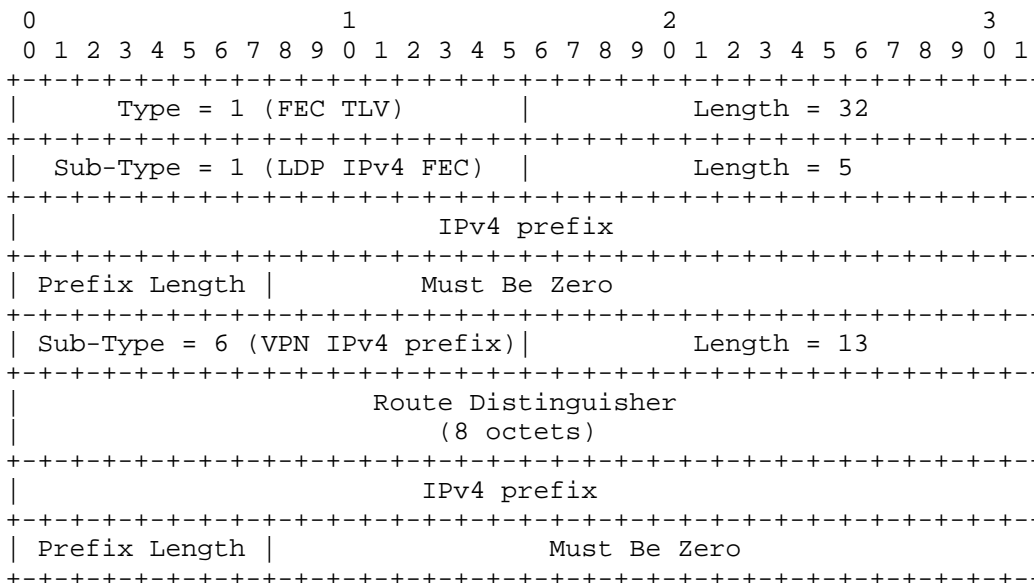


Types are defined below; Length is the length of the Value field in octets. The Value field depends on the Type; it is zero padded to align to a 4-octet boundary. TLVs may be nested within other TLVs, in which case the nested TLVs are called sub-TLVs. Sub-TLVs have independent types and MUST also be 4-octet aligned.

Two examples of how TLV and sub-TLV lengths are computed, and how sub-TLVs are padded to be 4-octet aligned, are as follows:



The Length for this TLV is 5. A Target FEC Stack TLV that contains an LDP IPv4 FEC sub-TLV and a VPN IPv4 prefix sub-TLV has the following format:



A description of the Types and Values of the top-level TLVs for LSP ping are given below:

Type #	Value Field
1	Target FEC Stack
2	Downstream Mapping (Deprecated)
3	Pad
4	Unassigned
5	Vendor Enterprise Number
6	Unassigned
7	Interface and Label Stack
8	Unassigned
9	Errored TLVs
10	Reply TOS Byte
20	Downstream Detailed Mapping

Types less than 32768 (i.e., with the high-order bit equal to 0) are mandatory TLVs that MUST either be supported by an implementation or result in the Return Code of 2 ("One or more of the TLVs was not understood") being sent in the echo response.

Types greater than or equal to 32768 (i.e., with the high-order bit equal to 1) are optional TLVs that SHOULD be ignored if the implementation does not understand or support them.

In Sections 3.2 through 3.9 and their various subsections, only the Value field of the TLV is included.

3.1. Return Codes

The Return Code is set to zero by the sender of an echo request. The receiver of said echo request can set it to one of the values listed below in the corresponding echo reply that it generates. The notation <RSC> refers to the Return Subcode. This field is filled in with the stack-depth for those codes that specify that. For all other codes, the Return Subcode MUST be set to zero.

Value	Meaning
-----	-----
0	No Return Code
1	Malformed echo request received
2	One or more of the TLVs was not understood
3	Replying router is an egress for the FEC at stack-depth <RSC>
4	Replying router has no mapping for the FEC at stack-depth <RSC>
5	Downstream Mapping Mismatch (See Note 1)
6	Upstream Interface Index Unknown (See Note 1)
7	Reserved
8	Label switched at stack-depth <RSC>
9	Label switched but no MPLS forwarding at stack-depth <RSC>
10	Mapping for this FEC is not the given label at stack-depth <RSC>
11	No label entry at stack-depth <RSC>
12	Protocol not associated with interface at FEC stack-depth <RSC>
13	Premature termination of ping due to label stack shrinking to a single label
14	See DDMAP TLV for meaning of Return Code and Return Subcode (See Note 2)
15	Label switched with FEC change

Note 1

The Return Subcode (RSC) contains the point in the label stack where processing was terminated. If the RSC is 0, no labels were processed. Otherwise, the packet was label switched at depth RSC.

Note 2

The Return Code is per "Downstream Detailed Mapping TLV" (Section 3.4). This Return Code MUST be used only in the message header and MUST be set only in the MPLS echo reply message. If the Return Code is set in the MPLS echo request message, then it MUST be ignored. When this Return Code is set, each Downstream Detailed Mapping TLV MUST have an appropriate Return Code and Return Subcode. This Return Code MUST be used when there are multiple downstreams for a given node (such as Point-to-Multipoint (P2MP) or ECMP), and the node needs to return a Return Code/Return Subcode for each downstream. This Return Code MAY be used even when there is only one downstream for a given node.

3.2. Target FEC Stack

A Target FEC Stack is a list of sub-TLVs. The number of elements is determined by looking at the sub-TLV length fields.

Sub-Type	Length	Value Field
-----	-----	-----
1	5	LDP IPv4 prefix
2	17	LDP IPv6 prefix
3	20	RSVP IPv4 LSP
4	56	RSVP IPv6 LSP
5		Unassigned
6	13	VPN IPv4 prefix
7	25	VPN IPv6 prefix
8	14	L2 VPN endpoint
9	10	"FEC 128" Pseudowire - IPv4 (deprecated)
10	14	"FEC 128" Pseudowire - IPv4
11	16+	"FEC 129" Pseudowire - IPv4
12	5	BGP labeled IPv4 prefix
13	17	BGP labeled IPv6 prefix
14	5	Generic IPv4 prefix
15	17	Generic IPv6 prefix
16	4	Nil FEC
24	38	"FEC 128" Pseudowire - IPv6
25	40+	"FEC 129" Pseudowire - IPv6

Other FEC types have been defined and will be defined as needed.

Note that this TLV defines a stack of FECs, the first FEC element corresponding to the top of the label stack, etc.

An MPLS echo request MUST have a Target FEC Stack that describes the FEC Stack being tested. For example, if an LSR X has an LDP mapping [RFC5036] for 192.0.2.1 (say, label 1001), then to verify that label 1001 does indeed reach an egress LSR that announced this prefix via LDP, X can send an MPLS echo request with a FEC Stack TLV with one FEC in it, namely, of type LDP IPv4 prefix, with prefix 192.0.2.1/32, and send the echo request with a label of 1001.

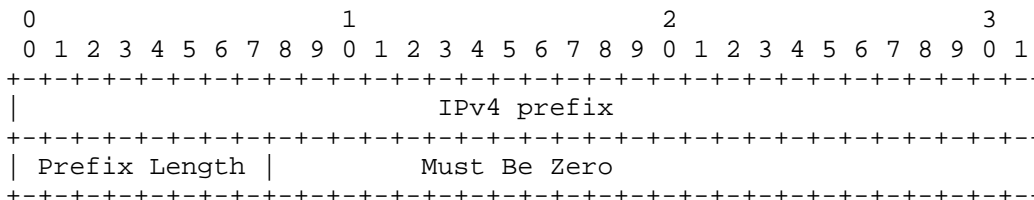
Say LSR X wanted to verify that a label stack of <1001, 23456> is the right label stack to use to reach a VPN IPv4 prefix (see Section 3.2.5) of 203.0.113.0/24 in VPN foo. Say further that LSR Y with loopback address 192.0.2.1 announced prefix 203.0.113.0/24 with Route Distinguisher (RD) RD-foo-Y (which may in general be different from the RD that LSR X uses in its own advertisements for VPN foo), label 23456, and BGP next hop 192.0.2.1 [RFC4271]. Finally, suppose that LSR X receives a label binding of 1001 for 192.0.2.1 via LDP. X has two choices in sending an MPLS echo request: X can send an MPLS echo request with a FEC Stack TLV with a single FEC of type VPN IPv4 prefix with a prefix of 203.0.113.0/24 and an RD of RD-foo-Y. Alternatively, X can send a FEC Stack TLV with two FECs, the first of type LDP IPv4 with a prefix of 192.0.2.1/32 and the second of type of IP VPN with a prefix 203.0.113.0/24 with an RD of RD-foo-Y. In either case, the MPLS echo request would have a label stack of <1001, 23456>. (Note: in this example, 1001 is the "outer" label and 23456 is the "inner" label.)

If, for example, an LSR Y has an LDP mapping for the IPv6 address 2001:db8::1 (say, label 2001), then to verify that label 2001 does reach an egress LSR that announced this prefix via LDP, LSR Y can send an MPLS echo request with a FEC Stack TLV with one LDP IPv6 prefix FEC, with prefix 2001:db8::1/128, and with a label of 2001.

If an end-to-end path comprises of one or more tunneled or stitched LSPs, each transit node that is the originating point of a new tunnel or segment SHOULD reply back notifying the FEC stack change along with the new FEC details, for example, if LSR X has an LDP mapping for IPv4 prefix 192.0.2.10 on LSR Z (say, label 3001). Say further that LSR A and LSR B are transit nodes along the path, which also have an RSVP tunnel over which LDP is enabled. While replying back, A SHOULD notify that the FEC changes from LDP to <RSVP, LDP>. If the new tunnel is a transparent pipe, i.e., the data-plane trace will not expire in the middle of the tunnel, then the transit node SHOULD NOT reply back notifying the FEC stack change or the new FEC details. If the transit node wishes to hide the nature of the tunnel from the ingress of the echo request, then the transit node MAY notify the FEC stack change and include Nil FEC as the new FEC.

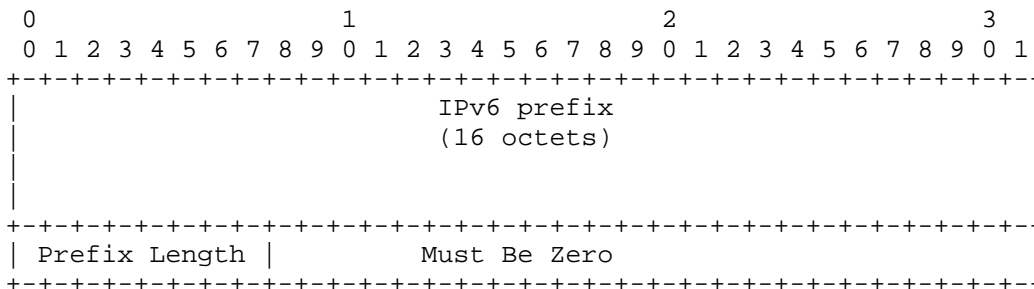
3.2.1. LDP IPv4 Prefix

The IPv4 Prefix FEC is defined in [RFC5036]. When an LDP IPv4 prefix is encoded in a label stack, the following format is used. The value consists of 4 octets of an IPv4 prefix followed by 1 octet of prefix length in bits; the format is given below. The IPv4 prefix is in network byte order; if the prefix is shorter than 32 bits, trailing bits SHOULD be set to zero. See [RFC5036] for an example of a Mapping for an IPv4 FEC.



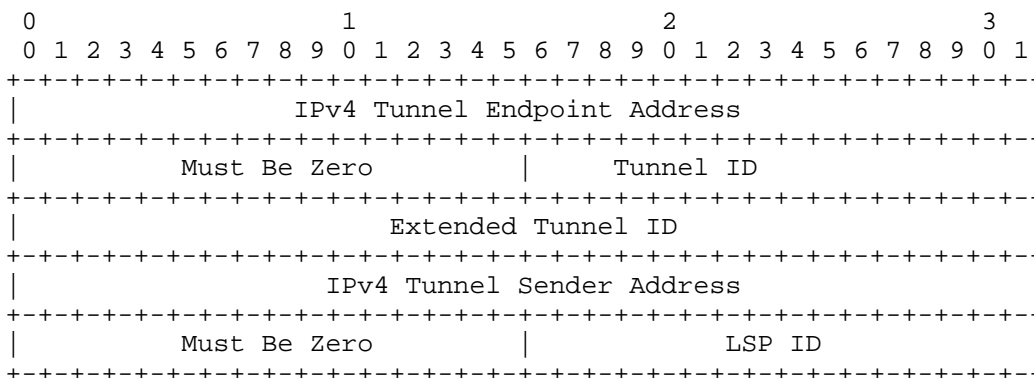
3.2.2. LDP IPv6 Prefix

The IPv6 Prefix FEC is defined in [RFC5036]. When an LDP IPv6 prefix is encoded in a label stack, the following format is used. The value consists of 16 octets of an IPv6 prefix followed by 1 octet of prefix length in bits; the format is given below. The IPv6 prefix is in network byte order; if the prefix is shorter than 128 bits, the trailing bits SHOULD be set to zero. See [RFC5036] for an example of a Mapping for an IPv6 FEC.



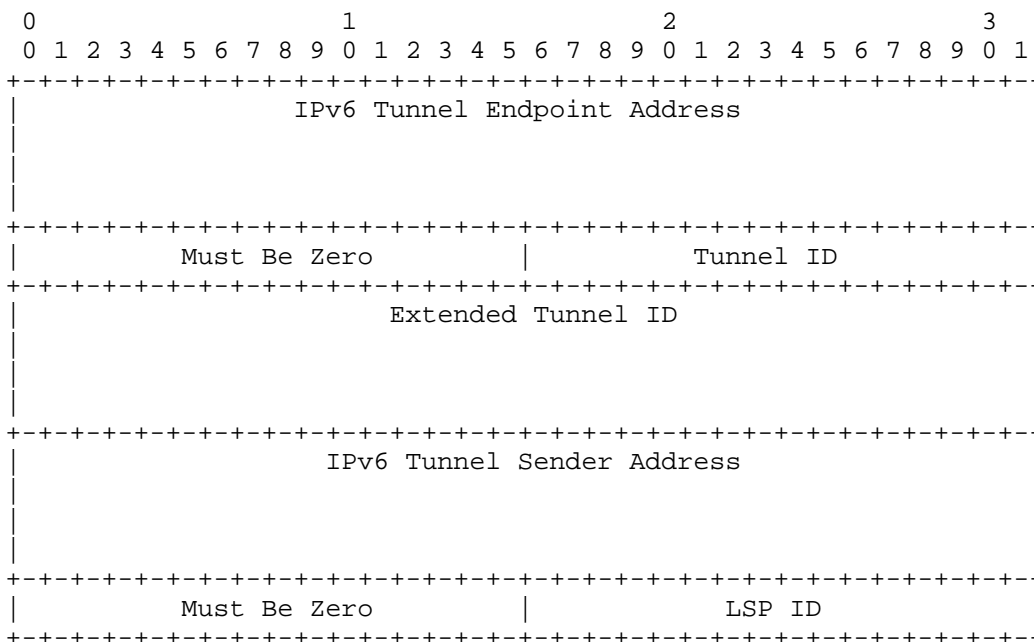
3.2.3. RSVP IPv4 LSP

The value has the format below. The Value fields are taken from RFC 3209 [RFC3209], Sections 4.6.1.1 and 4.6.2.1.



3.2.4. RSVP IPv6 LSP

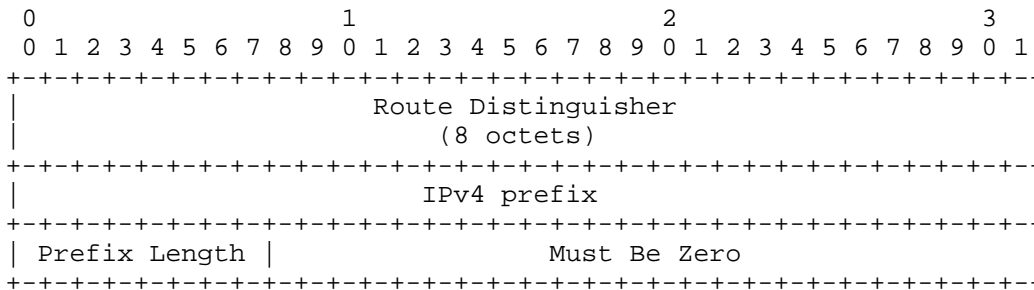
The value has the format below. The Value fields are taken from RFC 3209 [RFC3209], Sections 4.6.1.2 and 4.6.2.2.



3.2.5. VPN IPv4 Prefix

VPN-IPv4 Network Layer Routing Information (NLRI) is defined in [RFC4365]. This document uses the term VPN IPv4 prefix for a VPN-IPv4 NLRI that has been advertised with an MPLS label in BGP. See [RFC3107].

When a VPN IPv4 prefix is encoded in a label stack, the following format is used. The Value field consists of the RD advertised with the VPN IPv4 prefix, the IPv4 prefix (with trailing 0 bits to make 32 bits in all), and a prefix length, as follows:

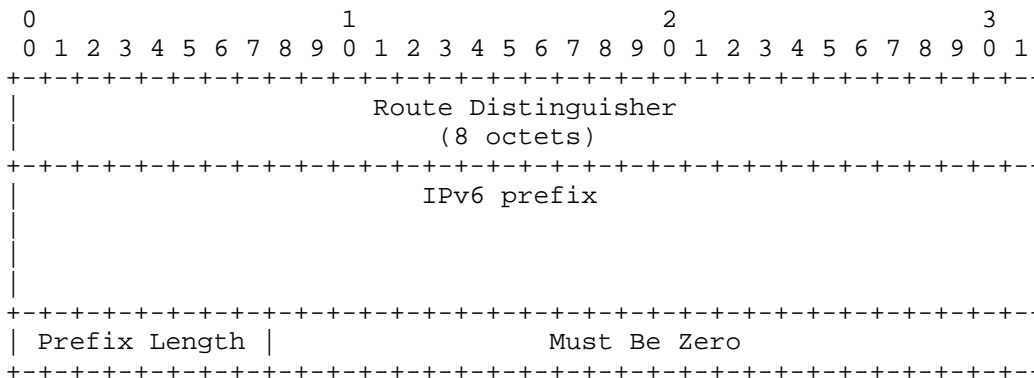


The RD is an 8-octet identifier; it does not contain any inherent information. The purpose of the RD is solely to allow one to create distinct routes to a common IPv4 address prefix. The encoding of the RD is not important here. When matching this field to the local FEC information, it is treated as an opaque value.

3.2.6. VPN IPv6 Prefix

VPN-IPv6 NLRI is defined in [RFC4365]. This document uses the term VPN IPv6 prefix for a VPN-IPv6 NLRI that has been advertised with an MPLS label in BGP. See [RFC3107].

When a VPN IPv6 prefix is encoded in a label stack, the following format is used. The Value field consists of the RD advertised with the VPN IPv6 prefix, the IPv6 prefix (with trailing 0 bits to make 128 bits in all), and a prefix length, as follows:

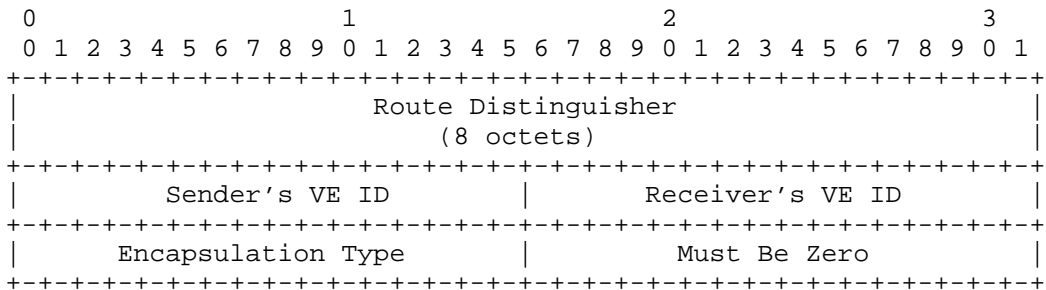


The RD is identical to the VPN IPv4 Prefix RD, except that it functions here to allow the creation of distinct routes to IPv6 prefixes. See Section 3.2.5. When matching this field to local FEC information, it is treated as an opaque value.

3.2.7. L2 VPN Endpoint

VPLS stands for Virtual Private LAN Service. The terms VPLS BGP NLRI and VPLS Edge Identifier (VE ID) are defined in [RFC4761]. This document uses the simpler term L2 VPN endpoint when referring to a VPLS BGP NLRI. The RD is an 8-octet identifier used to distinguish information about various L2 VPNs advertised by a node. The VE ID is a 2-octet identifier used to identify a particular node that serves as the service attachment point within a VPLS. The structure of these two identifiers is unimportant here; when matching these fields to local FEC information, they are treated as opaque values. The encapsulation type is identical to the Pseudowire (PW) Type in Section 3.2.9.

When an L2 VPN endpoint is encoded in a label stack, the following format is used. The Value field consists of an RD (8 octets), the sender's (of the ping) VE ID (2 octets), the receiver's VE ID (2 octets), and an encapsulation type (2 octets), formatted as follows:



3.2.8. FEC 128 Pseudowire - IPv4 (Deprecated)

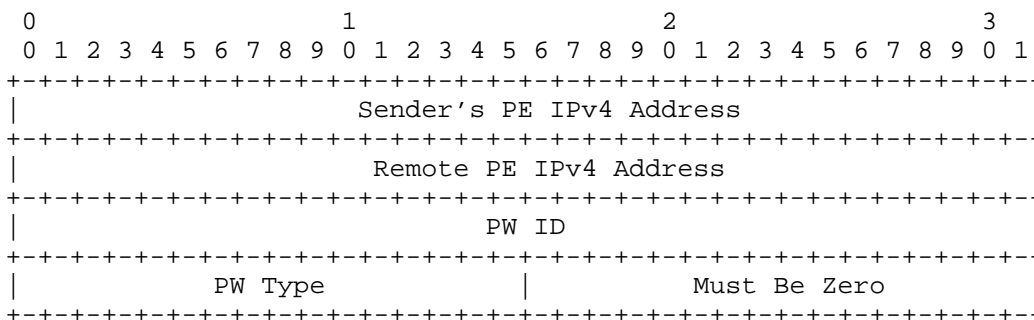
See Appendix A.1.1 for details.

3.2.9. FEC 128 Pseudowire - IPv4 (Current)

FEC 128 (0x80) is defined in [RFC8077], as are the terms PW ID (Pseudowire ID) and PW Type (Pseudowire Type). A PW ID is a non-zero 32-bit connection ID. The PW Type is a 15-bit number indicating the encapsulation type. It is carried right justified in the field below termed "encapsulation type" with the high-order bit set to zero.

Both of these fields are treated in this protocol as opaque values. When matching these fields to the local FEC information, the match MUST be exact.

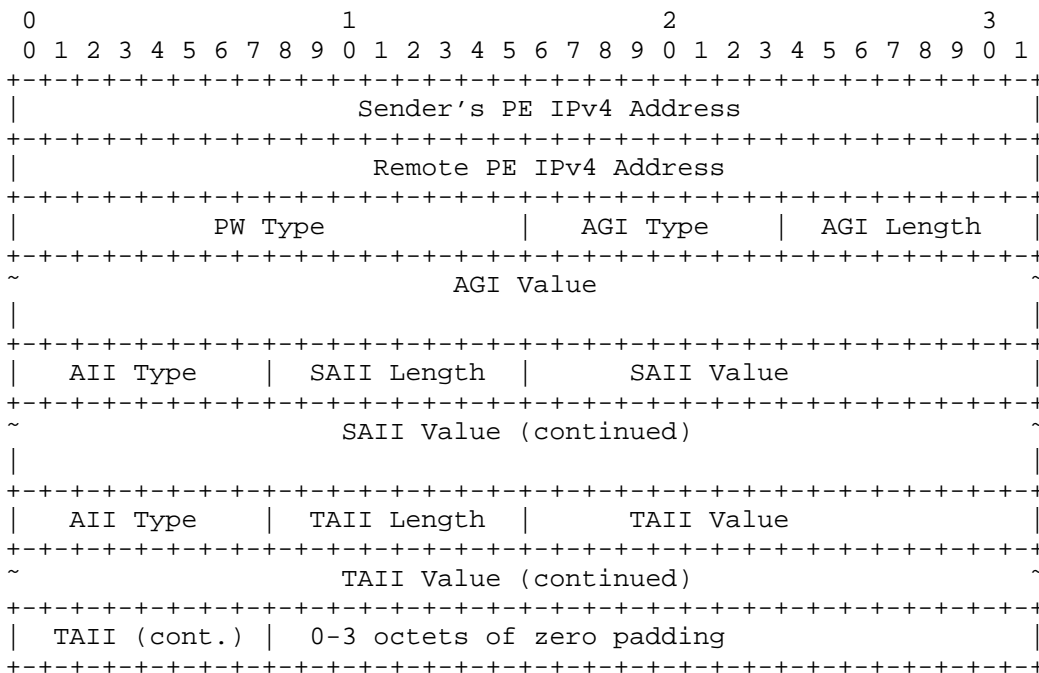
When a FEC 128 is encoded in a label stack, the following format is used. The Value field consists of the Sender's Provider Edge (PE) IPv4 Address (the source address of the targeted LDP session), the Remote PE IPv4 Address (the destination address of the targeted LDP session), the PW ID, and the encapsulation type as follows:



3.2.10. FEC 129 Pseudowire - IPv4

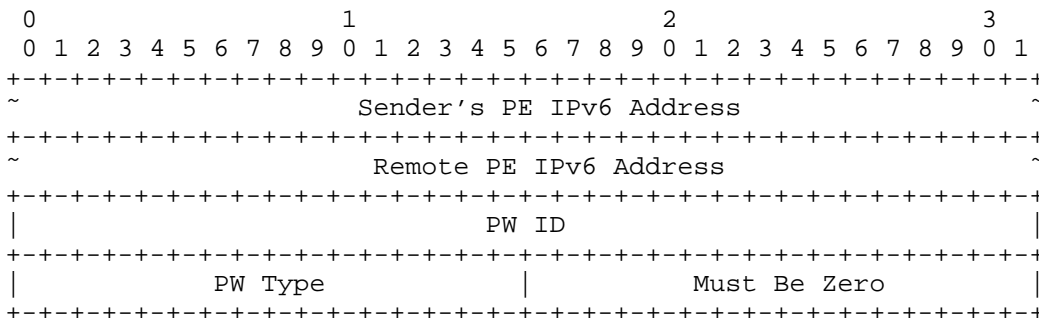
FEC 129 (0x81) and the terms PW Type, Attachment Group Identifier (AGI), Attachment Group Identifier Type (AGI Type), Attachment Individual Identifier Type (AII Type), Source Attachment Individual Identifier (SAII), and Target Attachment Individual Identifier (TAII) are defined in [RFC8077]. The PW Type is a 15-bit number indicating the encapsulation type. It is carried right justified in the field below PW Type with the high-order bit set to zero. All the other fields are treated as opaque values and copied directly from the FEC 129 format. All of these values together uniquely define the FEC within the scope of the LDP session identified by the source and remote PE IPv4 addresses.

When a FEC 129 is encoded in a label stack, the following format is used. The Length of this TLV is 16 + AGI length + SAII length + TAII length. Padding is used to make the total length a multiple of 4; the length of the padding is not included in the Length field.



3.2.11. FEC 128 Pseudowire - IPv6

The FEC 128 Pseudowire IPv6 sub-TLV has a structure consistent with the FEC 128 Pseudowire IPv4 sub-TLV as described in Section 3.2.9. The Value field consists of the Sender's PE IPv6 Address (the source address of the targeted LDP session), the Remote PE IPv6 Address (the destination address of the targeted LDP session), the PW ID, and the encapsulation type as follows:



Sender's PE IPv6 Address: The source IP address of the target IPv6 LDP session. 16 octets.

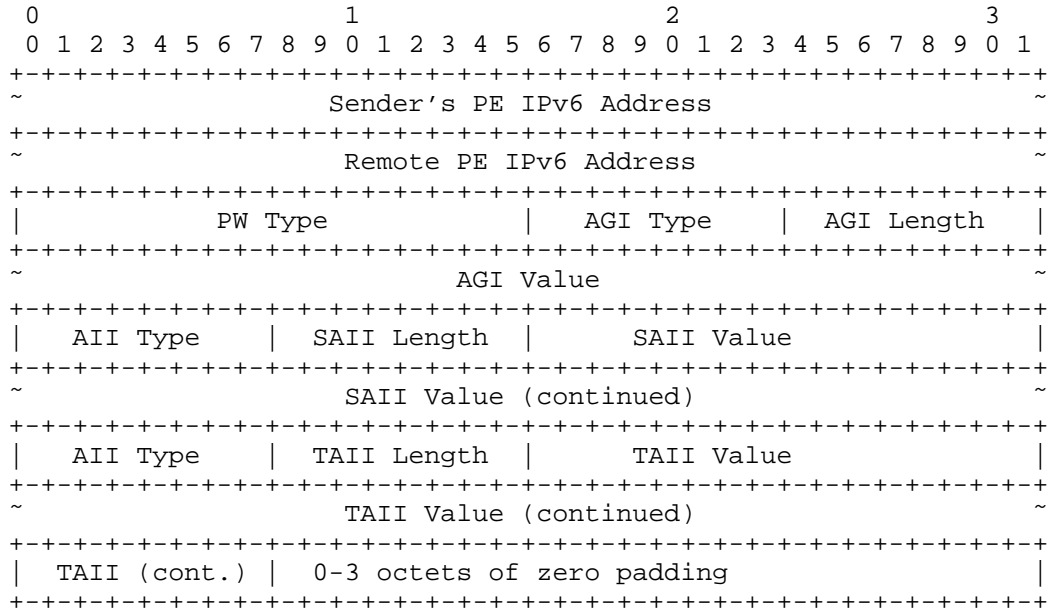
Remote PE IPv6 Address: The destination IP address of the target IPv6 LDP session. 16 octets.

PW ID: Same as FEC 128 Pseudowire IPv4 in Section 3.2.9.

PW Type: Same as FEC 128 Pseudowire IPv4 in Section 3.2.9.

3.2.12. FEC 129 Pseudowire - IPv6

The FEC 129 Pseudowire IPv6 sub-TLV has a structure consistent with the FEC 129 Pseudowire IPv4 sub-TLV as described in Section 3.2.10. When a FEC 129 is encoded in a label stack, the following format is used. The length of this TLV is 40 + AGI (Attachment Group Identifier) length + SAII (Source Attachment Individual Identifier) length + TAI (Target Attachment Individual Identifier) length. Padding is used to make the total length a multiple of 4; the length of the padding is not included in the Length field.



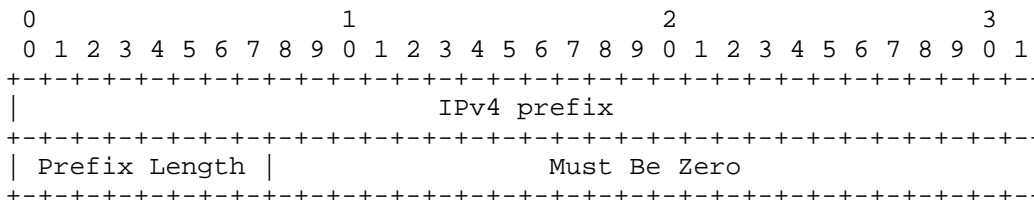
Sender's PE IPv6 Address: The source IP address of the target IPv6 LDP session. 16 octets.

Remote PE IPv6 Address: The destination IP address of the target IPv6 LDP session. 16 octets.

The other fields are the same as FEC 129 Pseudowire IPv4 in Section 3.2.10.

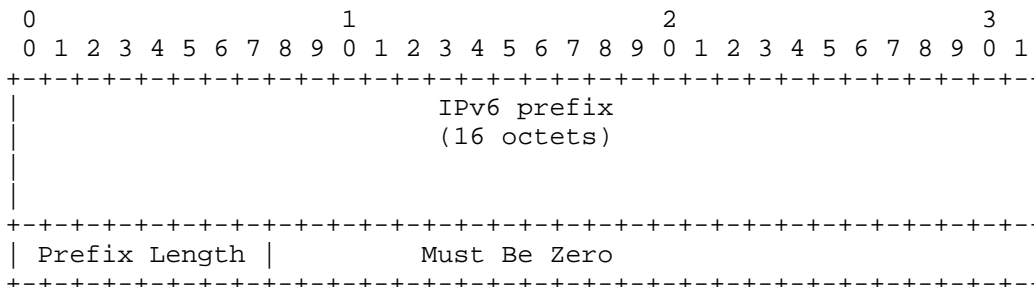
3.2.13. BGP Labeled IPv4 Prefix

BGP labeled IPv4 prefixes are defined in [RFC3107]. When a BGP labeled IPv4 prefix is encoded in a label stack, the following format is used. The Value field consists of the IPv4 prefix (with trailing 0 bits to make 32 bits in all) and the prefix length, as follows:



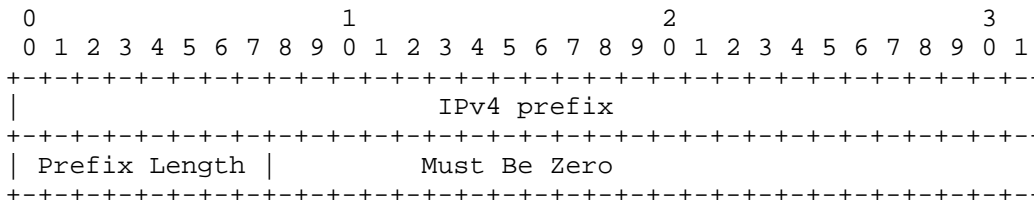
3.2.14. BGP Labeled IPv6 Prefix

BGP labeled IPv6 prefixes are defined in [RFC3107]. When a BGP labeled IPv6 prefix is encoded in a label stack, the following format is used. The value consists of 16 octets of an IPv6 prefix followed by 1 octet of prefix length in bits; the format is given below. The IPv6 prefix is in network byte order; if the prefix is shorter than 128 bits, the trailing bits SHOULD be set to zero.



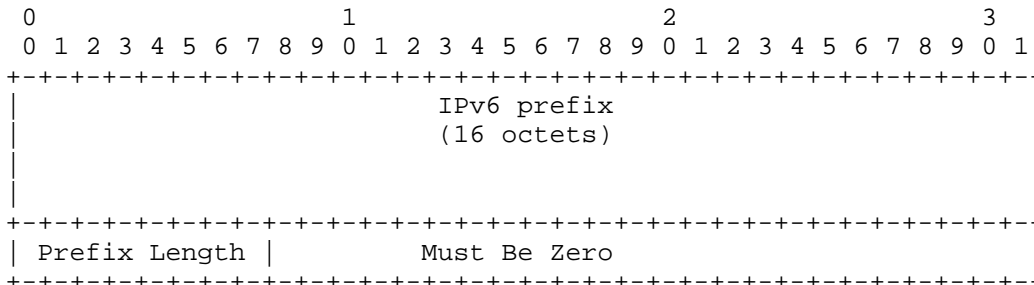
3.2.15. Generic IPv4 Prefix

The value consists of 4 octets of an IPv4 prefix followed by 1 octet of prefix length in bits; the format is given below. The IPv4 prefix is in network byte order; if the prefix is shorter than 32 bits, the trailing bits SHOULD be set to zero. This FEC is used if the protocol advertising the label is unknown or may change during the course of the LSP. An example is an inter-AS LSP that may be signaled by LDP in one Autonomous System (AS), by RSVP-TE [RFC3209] in another AS, and by BGP between the ASes, such as is common for inter-AS VPNs.



3.2.16. Generic IPv6 Prefix

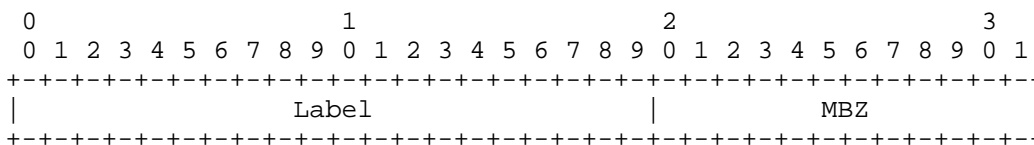
The value consists of 16 octets of an IPv6 prefix followed by 1 octet of prefix length in bits; the format is given below. The IPv6 prefix is in network byte order; if the prefix is shorter than 128 bits, the trailing bits SHOULD be set to zero.



3.2.17. Nil FEC

At times, labels from the reserved range, e.g., Router Alert and Explicit-null, may be added to the label stack for various diagnostic purposes such as influencing load-balancing. These labels may have no explicit FEC associated with them. The Nil FEC Stack is defined to allow a Target FEC Stack sub-TLV to be added to the Target FEC Stack to account for such labels so that proper validation can still be performed.

The Length is 4. Labels are 20-bit values treated as numbers.



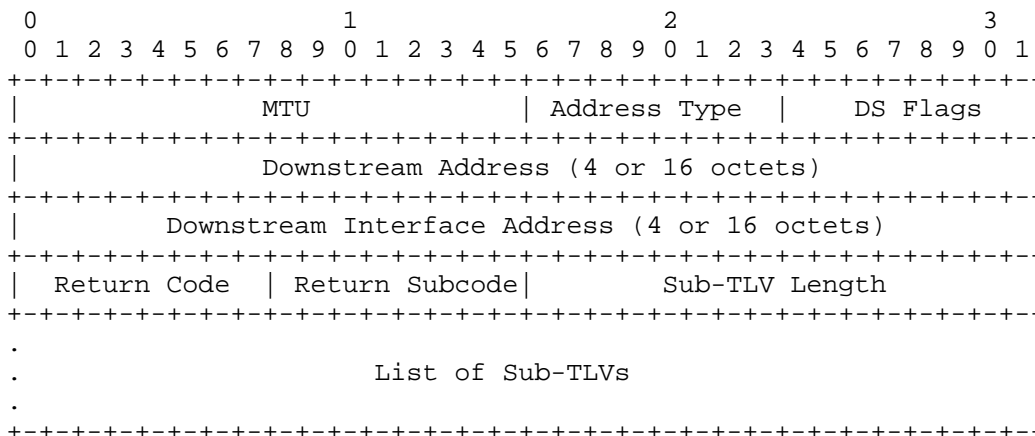
Label is the actual label value inserted in the label stack; the MBZ fields MUST be zero when sent and ignored on receipt.

3.3. Downstream Mapping (Deprecated)

See Appendix A.2 for more details.

3.4. Downstream Detailed Mapping TLV

The Downstream Detailed Mapping object is a TLV that MAY be included in an MPLS echo request message. Only one Downstream Detailed Mapping object may appear in an echo request. The presence of a Downstream Detailed Mapping object is a request that Downstream Detailed Mapping objects be included in the MPLS echo reply. If the replying router is the destination (Label Edge Router) of the FEC, then a Downstream Detailed Mapping TLV SHOULD NOT be included in the MPLS echo reply. Otherwise, the replying router SHOULD include a Downstream Detailed Mapping object for each interface over which this FEC could be forwarded. For a more precise definition of the notion of "downstream", see Section 3.4.2, "Downstream Router and Interface".



The Downstream Detailed Mapping TLV format is derived from the deprecated Downstream Mapping TLV format (see Appendix A.2.) The key change is that variable length and optional fields have been converted into sub-TLVs.

Maximum Transmission Unit (MTU)

The MTU is the size in octets of the largest MPLS frame (including label stack) that fits on the interface to the downstream LSR.

Address Type

The Address Type indicates if the interface is numbered or unnumbered. It also determines the length of the Downstream IP Address and Downstream Interface fields. The Address Type is set to one of the following values:

Type #	Address Type
-----	-----
1	IPv4 Numbered
2	IPv4 Unnumbered
3	IPv6 Numbered
4	IPv6 Unnumbered

DS Flags

The DS Flags field is a bit vector of various flags with the following format:

```

0 1 2 3 4 5 6 7
+-----+
| Rsvd(MBZ) |I|N|
+-----+
```

Two flags are defined currently, I and N. The remaining flags MUST be set to zero when sending and ignored on receipt.

Flag Name and Meaning

```

----
I Interface and Label Stack Object Request
```

When this flag is set, it indicates that the replying router SHOULD include an Interface and Label Stack Object in the echo reply message.

N Treat as a Non-IP Packet

Echo request messages will be used to diagnose non-IP flows. However, these messages are carried in IP packets. For a router that alters its ECMP algorithm based on the FEC or deep packet examination, this flag requests that the router treat this as it would if the determination of an IP payload had failed.

Downstream Address and Downstream Interface Address

IPv4 addresses and interface indices are encoded in 4 octets; IPv6 addresses are encoded in 16 octets.

If the interface to the downstream LSR is numbered, then the Address Type MUST be set to IPv4 or IPv6, the Downstream Address MUST be set to either the downstream LSR's Router ID or the interface address of the downstream LSR, and the Downstream Interface Address MUST be set to the downstream LSR's interface address.

If the interface to the downstream LSR is unnumbered, the Address Type MUST be IPv4 Unnumbered or IPv6 Unnumbered, the Downstream Address MUST be the downstream LSR's Router ID, and the Downstream Interface Address MUST be set to the index assigned by the upstream LSR to the interface.

If an LSR does not know the IP address of its neighbor, then it MUST set the Address Type to either IPv4 Unnumbered or IPv6 Unnumbered. For IPv4, it must set the Downstream Address to 127.0.0.1; for IPv6, the address is set to 0::1. In both cases, the interface index MUST be set to 0. If an LSR receives an Echo Request packet with either of these addresses in the Downstream Address field, this indicates that it MUST bypass interface verification but continue with label validation.

If the originator of an echo request packet wishes to obtain Downstream Detailed Mapping information but does not know the expected label stack, then it SHOULD set the Address Type to either IPv4 Unnumbered or IPv6 Unnumbered. For IPv4, it MUST set the Downstream Address to 224.0.0.2; for IPv6, the address MUST be set to FF02::2. In both cases, the interface index MUST be set to 0. If an LSR receives an echo request packet with the all-routers multicast address, then this indicates that it MUST bypass both interface and label stack validation but return Downstream Mapping TLVs using the information provided.

Return Code

The Return Code is set to zero by the sender of an echo request. The receiver of said echo request can set it in the corresponding echo reply that it generates to one of the values specified in Section 3.1 other than 14.

If the receiver sets a non-zero value of the Return Code field in the Downstream Detailed Mapping TLV, then the receiver MUST also set the Return Code field in the echo reply header to "See DDMAP TLV for Return Code and Return Subcode" (Section 3.1). An exception to this is if the receiver is a bud node [RFC4461] and is replying as both an egress and a transit node with a Return Code of 3 ("Replying router is an egress for the FEC at stack-depth <RSC>") in the echo reply header.

If the Return Code of the echo reply message is not set to either "See DDMAP TLV for Return Code and Return Subcode" (Section 3.1) or "Replying router is an egress for the FEC at stack-depth <RSC>", then the Return Code specified in the Downstream Detailed Mapping TLV MUST be ignored.

Return Subcode

The Return Subcode is set to zero by the sender. The receiver can set this field to an appropriate value as specified in Section 3.1: The Return Subcode is filled in with the stack-depth for those codes that specify the stack-depth. For all other codes, the Return Subcode MUST be set to zero.

If the Return Code of the echo reply message is not set to either "See DDMAP TLV for Return Code and Return Subcode" (Section 3.1) or "Replying router is an egress for the FEC at stack-depth <RSC>", then the Return Subcode specified in the Downstream Detailed Mapping TLV MUST be ignored.

Sub-TLV Length

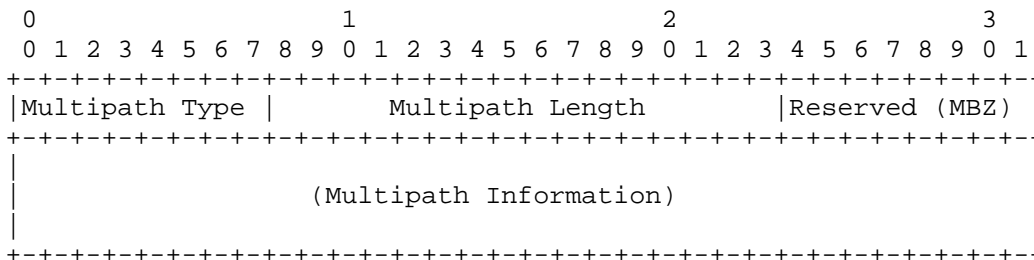
Total length in octets of the sub-TLVs associated with this TLV.

3.4.1. Sub-TLVs

This section defines the sub-TLVs that MAY be included as part of the Downstream Detailed Mapping TLV.

Sub-Type	Value Field
1	Multipath data
2	Label stack
3	FEC stack change

3.4.1.1. Multipath Data Sub-TLV



The multipath data sub-TLV includes Multipath Information.

Multipath Type

The type of the encoding for the Multipath Information.

The following Multipath Types are defined in this document:

Key	Type	Multipath Information
0	no multipath	Empty (Multipath Length = 0)
2	IP address	IP addresses
4	IP address range	low/high address pairs
8	Bit-masked IP address set	IP address prefix and bit mask
9	Bit-masked label set	Label prefix and bit mask

Type 0 indicates that all packets will be forwarded out this one interface.

Types 2, 4, 8, and 9 specify that the supplied Multipath Information will serve to exercise this path.

Multipath Length

The length in octets of the Multipath Information.

MBZ

MUST be set to zero when sending; MUST be ignored on receipt.

Multipath Information

Encoded multipath data (e.g., encoded address or label values), according to the Multipath Type. See Section 3.4.1.1.1 for encoding details.

3.4.1.1.1. Multipath Information Encoding

The Multipath Information encodes labels or addresses that will exercise this path. The Multipath Information depends on the Multipath Type. The contents of the field are shown in the table above. IPv4 addresses are drawn from the range 127/8; IPv6 addresses are drawn from the range 0:0:0:0:0:FFFF:7F00:0/104. Labels are treated as numbers, i.e., they are right justified in the field. For Type 4, ranges indicated by address pairs MUST NOT overlap and MUST be in ascending sequence.

Type 8 allows a more dense encoding of IP addresses. The IP prefix is formatted as a base IP address with the non-prefix low-order bits set to zero. The maximum prefix length is 27. Following the prefix is a mask of length 2^(32 - prefix length) bits for IPv4 and 2^(128 - prefix length) bits for IPv6. Each bit set to 1 represents a valid address. The address is the base IPv4 address plus the position of the bit in the mask where the bits are numbered left to right beginning with zero. For example, the IPv4 addresses 127.2.1.0, 127.2.1.5-127.2.1.15, and 127.2.1.20-127.2.1.29 would be encoded as follows:

0									1									2									3																						
0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9	0	1	2	3	4	5	6	7	8	9
+-----+																																																	
0 1 1 1 1 1 1 1 0 0 0 0 0 0 1 0 0 0 0 0 0 0 0 0 1 0 0 0 0 0 0 0																																																	
+-----+																																																	
1 0 0 0 0 1 1 1 1 1 1 1 1 1 1 1 0 0 0 0 1 1 1 1 1 1 1 1 1 1 1 0 0																																																	
+-----+																																																	

Those same addresses embedded in IPv6 would be encoded as follows:

```

      0                1                2                3
0 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+-----+-----+-----+-----+-----+-----+-----+-----+
|0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0|
+-----+-----+-----+-----+-----+-----+-----+-----+
|0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0|
+-----+-----+-----+-----+-----+-----+-----+-----+
|0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1|
+-----+-----+-----+-----+-----+-----+-----+-----+
|0 1 1 1 1 1 1 1 0 0 0 0 0 0 1 0 0 0 0 0 0 0 0 0 1 0 0 0 0 0 0 0|
+-----+-----+-----+-----+-----+-----+-----+-----+
|1 0 0 0 0 1 1 1 1 1 1 1 1 1 1 1 0 0 0 0 1 1 1 1 1 1 1 1 1 1 0 0|
+-----+-----+-----+-----+-----+-----+-----+-----+

```

Type 9 allows a more dense encoding of labels. The label prefix is formatted as a base label value with the non-prefix low-order bits set to zero. The maximum prefix (including leading zeros due to encoding) length is 27. Following the prefix is a mask of length $2^{(32 - \text{prefix length})}$ bits. Each bit set to one represents a valid label. The label is the base label plus the position of the bit in the mask where the bits are numbered left to right beginning with zero. Label values of all the odd numbers between 1152 and 1279 would be encoded as follows:

```

      0                1                2                3
0 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
+-----+-----+-----+-----+-----+-----+-----+-----+
|0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 1 0 0 1 0 0 0 0 0 0|
+-----+-----+-----+-----+-----+-----+-----+-----+
|0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1|
+-----+-----+-----+-----+-----+-----+-----+-----+
|0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1|
+-----+-----+-----+-----+-----+-----+-----+-----+
|0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1 0 1|
+-----+-----+-----+-----+-----+-----+-----+-----+

```

If the received Multipath Information is non-null, the labels and IP addresses MUST be picked from the set provided. If none of these labels or addresses map to a particular downstream interface, then for that interface, the type MUST be set to 0. If the received Multipath Information is null (i.e., Multipath Length = 0, or for Types 8 and 9, a mask of all zeros), the type MUST be set to 0.

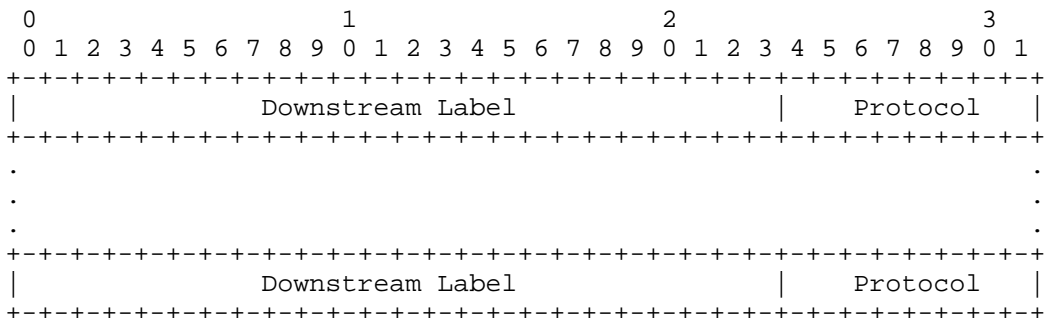
For example, suppose LSR X at hop 10 has two downstream LSRs, Y and Z, for the FEC in question. The received X could return Multipath Type 4, with low/high IP addresses of 127.1.1.1->127.1.1.255 for downstream LSR Y and 127.2.1.1->127.2.1.255 for downstream LSR Z. The head end reflects this information to LSR Y. Y, which has three downstream LSRs, U, V, and W, computes that 127.1.1.1->127.1.1.127 would go to U and 127.1.1.128-> 127.1.1.255 would go to V. Y would then respond with 3 Downstream Detailed Mapping TLVs: to U, with Multipath Type 4 (127.1.1.1->127.1.1.127); to V, with Multipath Type 4 (127.1.1.127->127.1.1.255); and to W, with Multipath Type 0.

Note that computing Multipath Information may impose a significant processing burden on the receiver. A receiver MAY thus choose to process a subset of the received prefixes. The sender, on receiving a reply to a Downstream Detailed Mapping with partial information, SHOULD assume that the prefixes missing in the reply were skipped by the receiver and MAY re-request information about them in a new echo request.

The encoding of Multipath Information in scenarios where a few LSRs apply Entropy-label-based load-balancing while other LSRs are non-EL (IP-based) load balanced will be defined in a different document.

The encoding of Multipath Information in scenarios where LSRs have Layer 2 ECMP over Link Aggregation Group (LAG) interfaces will be defined in a different document.

3.4.1.2. Label Stack Sub-TLV



The Label Stack sub-TLV contains the set of labels in the label stack as it would have appeared if this router were forwarding the packet through this interface. Any Implicit Null labels are explicitly included. The number of label/protocol pairs present in the sub-TLV is determined based on the sub-TLV data length. When the Downstream Detailed Mapping TLV is sent in the echo reply, this sub-TLV MUST be included.

Downstream Label

A downstream label is 24 bits, in the same format as an MPLS label minus the TTL field, i.e., the MSBit of the label is bit 0, the LSBit is bit 19, the TC field [RFC5462] is bits 20-22, and S is bit 23. The replying router SHOULD fill in the TC field and S bit; the LSR receiving the echo reply MAY choose to ignore these.

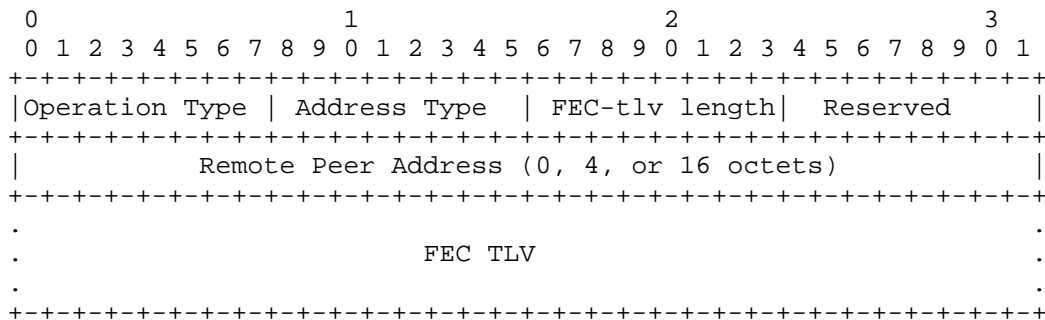
Protocol

This specifies the label distribution protocol for the Downstream label. Protocol values are taken from the following table:

Protocol #	Signaling Protocol
0	Unknown
1	Static
2	BGP
3	LDP
4	RSVP-TE

3.4.1.3. FEC Stack Change Sub-TLV

A router MUST include the FEC stack change sub-TLV when the downstream node in the echo reply has a different FEC Stack than the FEC Stack received in the echo request. One or more FEC stack change sub-TLVs MAY be present in the Downstream Detailed Mapping TLV. The format is as below.



Operation Type

The operation type specifies the action associated with the FEC stack change. The following operation types are defined:

Type #	Operation
-----	-----
1	Push
2	Pop

Address Type

The Address Type indicates the remote peer's address type. The Address Type is set to one of the following values. The length of the peer address is determined based on the address type. The address type MAY be different from the address type included in the Downstream Detailed Mapping TLV. This can happen when the LSP goes over a tunnel of a different address family. The address type MAY be set to Unspecified if the peer address is either unavailable or the transit router does not wish to provide it for security or administrative reasons.

Type #	Address Type	Address length
-----	-----	-----
0	Unspecified	0
1	IPv4	4
2	IPv6	16

FEC TLV Length

Length in octets of the FEC TLV.

Reserved

This field is reserved for future use and MUST be set to zero.

Remote Peer Address

The remote peer address specifies the remote peer that is the next hop for the FEC being currently traced. If the operation type is PUSH, the remote peer address is the address of the peer from which the FEC being pushed was learned. If the operation type is pop, the remote peer address MAY be set to Unspecified.

For upstream-assigned labels [RFC5331], an operation type of pop will have a remote peer address (the upstream node that assigned the label), and this SHOULD be included in the FEC stack change

sub-TLV. The remote peer address MAY be set to Unspecified if the address needs to be hidden.

FEC TLV

The FEC TLV is present only when the FEC-tlv length field is non-zero. The FEC TLV specifies the FEC associated with the FEC stack change operation. This TLV MAY be included when the operation type is pop. It MUST be included when the operation type is PUSH. The FEC TLV contains exactly one FEC from the list of FECs specified in Section 3.2. A Nil FEC MAY be associated with a PUSH operation if the responding router wishes to hide the details of the FEC being pushed.

FEC stack change sub-TLV operation rules are as follows:

- a. A FEC stack change sub-TLV containing a PUSH operation MUST NOT be followed by a FEC stack change sub-TLV containing a pop operation.
- b. One or more pop operations MAY be followed by one or more PUSH operations.
- c. One FEC stack change sub-TLV MUST be included per FEC stack change. For example, if 2 labels are going to be pushed, then one FEC stack change sub-TLV MUST be included for each FEC.
- d. A FEC splice operation (an operation where one FEC ends and another FEC starts, MUST be performed by including a pop type FEC stack change sub-TLV followed by a PUSH type FEC stack change sub-TLV.
- e. A Downstream Detailed Mapping TLV containing only one FEC stack change sub-TLV with pop operation is equivalent to IS_EGRESS (Return Code 3, Section 3.1) for the outermost FEC in the FEC stack. The ingress router performing the LSP traceroute MUST treat such a case as an IS_EGRESS for the outermost FEC.

3.4.2. Downstream Router and Interface

The notion of "downstream router" and "downstream interface" should be explained. Consider an LSR X. If a packet that was originated with TTL $n > 1$ arrived with outermost label L and TTL=1 at LSR X, X must be able to compute which LSRs could receive the packet if it was originated with TTL= $n+1$, over which interface the request would arrive and what label stack those LSRs would see. (It is outside the scope of this document to specify how this computation is done.) The set of these LSRs/interfaces consists of the downstream routers/

interfaces (and their corresponding labels) for X with respect to L. Each pair of downstream router and interface requires a separate Downstream Detailed Mapping to be added to the reply.

The case where X is the LSR originating the echo request is a special case. X needs to figure out what LSRs would receive the MPLS echo request for a given FEC Stack that X originates with TTL=1.

The set of downstream routers at X may be alternative paths (see the discussion below on ECMP) or simultaneous paths (e.g., for MPLS multicast). In the former case, the Multipath Information is used as a hint to the sender as to how it may influence the choice of these alternatives.

3.5. Pad TLV

The value part of the Pad TLV contains a variable number (≥ 1) of octets. The first octet takes values from the following table; all the other octets (if any) are ignored. The receiver SHOULD verify that the TLV is received in its entirety, but otherwise ignores the contents of this TLV, apart from the first octet.

Value	Meaning
-----	-----
0	Reserved
1	Drop Pad TLV from reply
2	Copy Pad TLV to reply
3-250	Unassigned
251-254	Reserved for Experimental Use
255	Reserved

The Pad TLV can be added to an echo request to create a message of a specific length in cases where messages of various sizes are needed for troubleshooting. The first octet allows for controlling the inclusion of this additional padding in the respective echo reply.

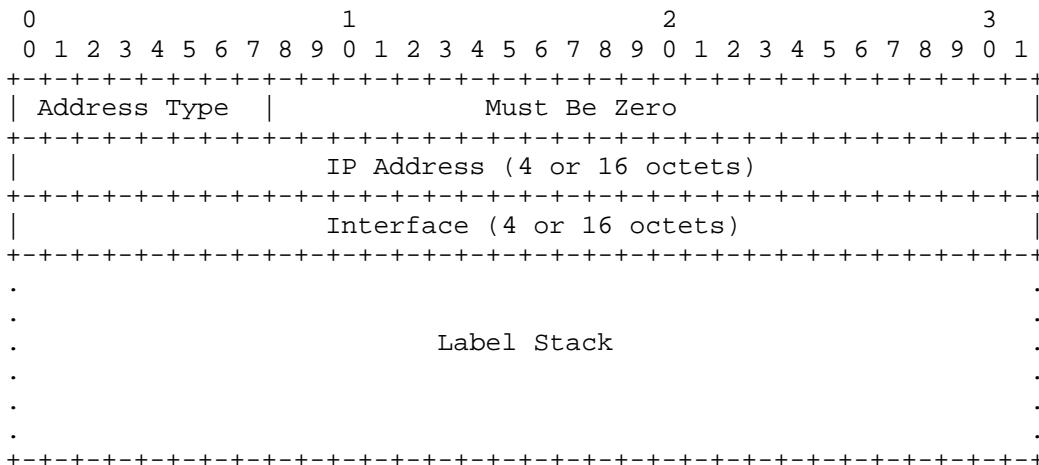
3.6. Vendor Enterprise Number

"Private Enterprise Numbers" [IANA-ENT] are maintained by IANA. The Length of this TLV is always 4; the value is the Structure of Management Information (SMI) Private Enterprise Code, in network octet order, of the vendor with a Vendor Private extension to any of the fields in the fixed part of the message, in which case this TLV MUST be present. If none of the fields in the fixed part of the message have Vendor Private extensions, inclusion of this TLV is OPTIONAL. Vendor Private ranges for Message Types, Reply Modes, and Return Codes have been defined. When any of these are used, the Vendor Enterprise Number TLV MUST be included in the message.

3.7. Interface and Label Stack

The Interface and Label Stack TLV MAY be included in a reply message to report the interface on which the request message was received and the label stack that was on the packet when it was received. Only one such object may appear. The purpose of the object is to allow the upstream router to obtain the exact interface and label stack information as it appears at the replying LSR.

The Length is $K + 4*N$ octets; N is the number of labels in the label stack. Values for K are found in the description of Address Type below. The Value field of this TLV has the following format:



Address Type

The Address Type indicates if the interface is numbered or unnumbered. It also determines the length of the IP Address and Interface fields. The resulting total for the initial part of the TLV is listed in the table below as "K Octets". The Address Type is set to one of the following values:

Type #	Address Type	K Octets
-----	-----	-----
0	Reserved	4
1	IPv4 Numbered	12
2	IPv4 Unnumbered	12
3	IPv6 Numbered	36
4	IPv6 Unnumbered	24
5-250	Unassigned	
251-254	Reserved for Experimental Use	
255	Reserved	

IP Address and Interface

IPv4 addresses and interface indices are encoded in 4 octets; IPv6 addresses are encoded in 16 octets.

If the interface upon which the echo request message was received is numbered, then the Address Type MUST be set to IPv4 or IPv6, the IP Address MUST be set to either the LSR's Router ID or the interface address, and the Interface MUST be set to the interface address.

If the interface is unnumbered, the Address Type MUST be either IPv4 Unnumbered or IPv6 Unnumbered, the IP Address MUST be the LSR's Router ID, and the Interface MUST be set to the index assigned to the interface.

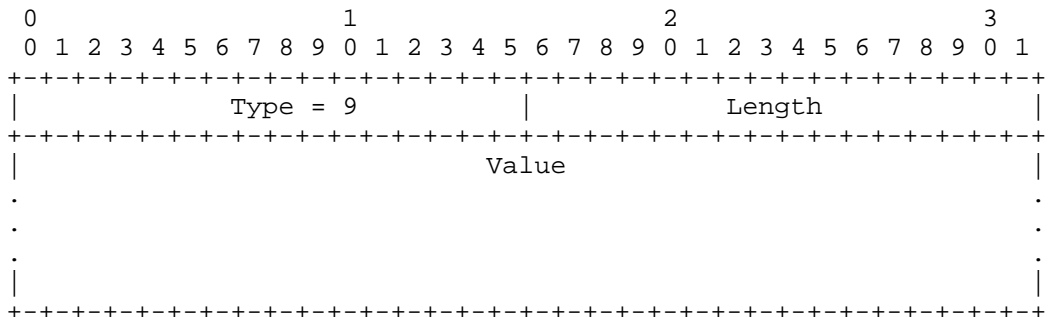
Label Stack

The label stack of the received echo request message. If any TTL values have been changed by this router, they SHOULD be restored.

3.8. Errored TLVs

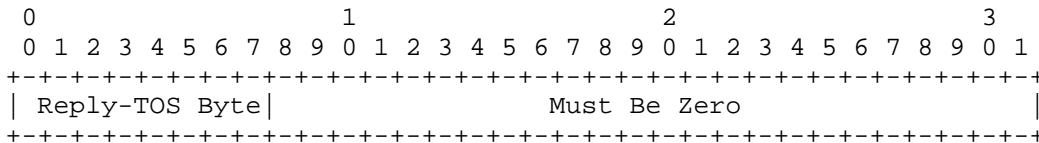
The following TLV is a TLV that MAY be included in an echo reply to inform the sender of an echo request of mandatory TLVs either not supported by an implementation or parsed and found to be in error.

The Value field contains the TLVs that were not understood, encoded as sub-TLVs.



3.9. Reply TOS Octet TLV

This TLV MAY be used by the originator of the echo request to request that an echo reply be sent with the IP header Type of Service (TOS) octet set to the value specified in the TLV. This TLV has a length of 4 with the following Value field.



4. Theory of Operation

An MPLS echo request is used to test a particular LSP. The LSP to be tested is identified by the "FEC Stack"; for example, if the LSP was set up via LDP, and a label is mapped to an egress IP address of 198.51.100.1, the FEC Stack contains a single element, namely, an LDP IPv4 prefix sub-TLV with value 198.51.100.1/32. If the LSP being tested is an RSVP LSP, the FEC Stack consists of a single element that captures the RSVP Session and Sender Template that uniquely identifies the LSP.

FEC Stacks can be more complex. For example, one may wish to test a VPN IPv4 prefix of 203.0.113.0/24 that is tunneled over an LDP LSP with egress 192.0.2.1. The FEC Stack would then contain two sub-TLVs, the bottom being a VPN IPv4 prefix, and the top being an LDP IPv4 prefix. If the underlying (LDP) tunnel were not known, or was considered irrelevant, the FEC Stack could be a single element with just the VPN IPv4 sub-TLV.

When an MPLS echo request is received, the receiver is expected to verify that the control plane and data plane are both healthy (for the FEC Stack being pinged), and that the two planes are in sync. The procedures for this are in Section 4.4.

4.1. Dealing with Equal-Cost Multipath (ECMP)

LSPs need not be simple point-to-point tunnels. Frequently, a single LSP may originate at several ingresses and terminate at several egresses; this is very common with LDP LSPs. LSPs for a given FEC may also have multiple "next hops" at transit LSRs. At an ingress, there may also be several different LSPs to choose from to get to the desired endpoint. Finally, LSPs may have backup paths, detour paths, and other alternative paths to take should the primary LSP go down.

Regarding the last two points stated above: it is assumed that the LSR sourcing MPLS echo requests can force the echo request into any desired LSP, so choosing among multiple LSPs at the ingress is not an issue. The problem of probing the various flavors of backup paths that will typically not be used for forwarding data unless the primary LSP is down will not be addressed here.

Since the actual LSP and path that a given packet may take may not be known a priori, it is useful if MPLS echo requests can exercise all possible paths. This, although desirable, may not be practical because the algorithms that a given LSR uses to distribute packets over alternative paths may be proprietary.

To achieve some degree of coverage of alternate paths, there is a certain latitude in choosing the destination IP address and source UDP port for an MPLS echo request. This is clearly not sufficient; in the case of traceroute, more latitude is offered by means of the Multipath Information of the Downstream Detailed Mapping TLV. This is used as follows. An ingress LSR periodically sends an LSP traceroute message to determine whether there are multipaths for a given LSP. If so, each hop will provide some information as to how each of its downstream paths can be exercised. The ingress can then send MPLS echo requests that exercise these paths. If several transit LSRs have ECMP, the ingress may attempt to compose these to exercise all possible paths. However, full coverage may not be possible.

4.2. Testing LSPs That Are Used to Carry MPLS Payloads

To detect certain LSP breakages, it may be necessary to encapsulate an MPLS echo request packet with at least one additional label when testing LSPs that are used to carry MPLS payloads (such as LSPs used to carry L2VPN and L3VPN traffic. For example, when testing LDP or RSVP-TE LSPs, just sending an MPLS echo request packet may not detect instances where the router immediately upstream of the destination of the LSP ping may forward the MPLS echo request successfully over an interface not configured to carry MPLS payloads because of the use of penultimate hop popping. Since the receiving router has no means to ascertain whether the IP packet was sent unlabeled or implicitly labeled, the addition of labels shimmed above the MPLS echo request (using the Nil FEC) will prevent a router from forwarding such a packet out to unlabeled interfaces.

4.3. Sending an MPLS Echo Request

An MPLS echo request is a UDP packet. The IP header is set as follows: the source IP address is a routable address of the sender; the destination IP address is a (randomly chosen) IPv4 address from the range 127/8 or an IPv6 address from the range 0:0:0:0:0:FFFF:7F00:0/104. The IP TTL is set to 1. The source UDP port is chosen by the sender; the destination UDP port is set to 3503 (assigned by IANA for MPLS echo requests). The Router Alert IP Option of value 0x0 [RFC2113] for IPv4 or value 69 [RFC7506] for IPv6 MUST be set in the IP header.

An MPLS echo request is sent with a label stack corresponding to the FEC Stack being tested. Note that further labels could be applied if, for example, the normal route to the topmost FEC in the stack is via a Traffic Engineered Tunnel [RFC3209]. If all of the FECs in the stack correspond to Implicit Null labels, the MPLS echo request is considered unlabeled even if further labels will be applied in sending the packet.

If the echo request is labeled, one MAY (depending on what is being pinged) set the TTL of the innermost label to 1, to prevent the ping request going farther than it should. Examples of where this SHOULD be done include pinging a VPN IPv4 or IPv6 prefix, an L2 VPN endpoint, or a pseudowire. Preventing the ping request from going too far can also be accomplished by inserting a Router Alert label above this label; however, this may lead to the undesired side effect that MPLS echo requests take a different data path than actual data. For more information on how these mechanisms can be used for pseudowire connectivity verification, see [RFC5085][RFC5885].

In "ping" mode (end-to-end connectivity check), the TTL in the outermost label is set to 255. In "traceroute" mode (fault isolation mode), the TTL is set successively to 1, 2, and so on.

The sender chooses a Sender's Handle and a Sequence Number. When sending subsequent MPLS echo requests, the sender SHOULD increment the Sequence Number by 1. However, a sender MAY choose to send a group of echo requests with the same Sequence Number to improve the chance of arrival of at least one packet with that Sequence Number.

The TimeStamp Sent is set to the time of day in NTP format that the echo request is sent. The TimeStamp Received is set to zero.

An MPLS echo request MUST have a FEC Stack TLV. Also, the Reply Mode must be set to the desired Reply Mode; the Return Code and Subcode are set to zero. In the "traceroute" mode, the echo request SHOULD include a Downstream Detailed Mapping TLV.

4.4. Receiving an MPLS Echo Request

Sending an MPLS echo request to the control plane is triggered by one of the following packet processing exceptions: Router Alert option, IP TTL expiration, MPLS TTL expiration, MPLS Router Alert label, or the destination address in the 127/8 address range. The control plane further identifies it by UDP destination port 3503.

For reporting purposes, the bottom of the stack is considered to be a stack-depth of 1. This is to establish an absolute reference for the case where the actual stack may have more labels than there are FECs in the Target FEC Stack.

Furthermore, in all the Return Codes listed in this document, a stack-depth of 0 means "no value specified". This allows compatibility with existing implementations that do not use the Return Subcode field.

An LSR X that receives an MPLS echo request then processes it as follows.

1. General packet sanity is verified. If the packet is not well-formed, LSR X SHOULD send an MPLS echo reply with the Return Code set to "Malformed echo request received" and the Subcode set to zero. If there are any TLVs not marked as "Ignore" (i.e., if the TLV type is less than 32768, see Section 3) that LSR X does not understand, LSR X SHOULD send an MPLS "TLV not understood" (as appropriate), and set the Subcode to zero. In the latter case, the misunderstood TLVs (only) are included as sub-TLVs in an Errored TLVs TLV in the reply. The header field's Sender's Handle, Sequence Number, and Timestamp Sent are not examined but are included in the MPLS echo reply message.

The algorithm uses the following variables and identifiers:

Interface-I:	the interface on which the MPLS echo request was received.
Stack-R:	the label stack on the packet as it was received.
Stack-D:	the label stack carried in the "Label stack sub-TLV" in the Downstream Detailed Mapping TLV (not always present).
Label-L:	the label from the actual stack currently being examined. Requires no initialization.

Label-stack-depth: the depth of the label being verified. Initialized to the number of labels in the received label stack S.

FEC-stack-depth: depth of the FEC in the Target FEC Stack that should be used to verify the current actual label. Requires no initialization.

Best-return-code: contains the Return Code for the echo reply packet as currently best known. As the algorithm progresses, this code may change depending on the results of further checks that it performs.

Best-rtn-subcode: similar to Best-return-code, but for the echo reply Subcode.

FEC-status: result value returned by the FEC Checking algorithm described in Section 4.4.1.

```
/* Save receive context information */
```

2. If the echo request is good, LSR X stores the interface over which the echo was received in Interface-I, and the label stack with which it came in Stack-R.

```
/* The rest of the algorithm iterates over the labels in Stack-R, verifies validity of label values, reports associated label switching operations (for traceroute), verifies correspondence between the Stack-R and the Target FEC Stack description in the body of the echo request, and reports any errors. */
```

```
/* The algorithm iterates as follows. */
```

3. Label Validation:

```
If Label-stack-depth is 0 {
```

```
/* The LSR needs to report that it is a tail end for the LSP */
```

```
Set FEC-stack-depth to 1, set Label-L to 3 (Implicit Null).  
Set Best-return-code to 3 ("Replying router is an egress for the FEC at stack-depth"), set Best-rtn-subcode to the value of FEC-stack-depth (1), and go to step 5 (Egress Processing).
```

```
}
```

```
/* This step assumes there is always an entry for well-known label values */
```


Set Label-L to the value extracted from Stack-R at depth Label-stack-depth. Look up Label-L in the Incoming Label Map (ILM) to determine if the label has been allocated and an operation is associated with it.

If there is no entry for Label-L {

/* Indicates a temporary or permanent label synchronization problem, and the LSR needs to report an error */

Set Best-return-code to 11 ("No label entry at stack-depth") and Best-rtn-subcode to Label-stack-depth. Go to step 7 (Send Reply Packet).

}

Else {

Retrieve the associated label operation from the corresponding Next Hop Label Forwarding Entry (NHLFE), and proceed to step 4 (Label Operation Check).

}

4. Label Operation Check

If the label operation is "Pop and Continue Processing" {

/* Includes Explicit Null and Router Alert label cases */

Iterate to the next label by decrementing Label-stack-depth, and loop back to step 3 (Label Validation).

}

If the label operation is "Swap or Pop and Switch based on Popped Label" {

Set Best-return-code to 8 ("Label switched at stack-depth") and Best-rtn-subcode to Label-stack-depth to report transit switching.

If a Downstream Detailed Mapping TLV is present in the received echo request {

If the IP address in the TLV is 127.0.0.1 or 0::1 {

```
Set Best-return-code to 6 ("Upstream Interface Index
Unknown"). An Interface and Label Stack TLV SHOULD be
included in the reply and filled with Interface-I and
Stack-R.

}

Else {

    Verify that the IP address, interface address, and label
    stack in the Downstream Detailed Mapping TLV match
    Interface-I and Stack-R. If there is a mismatch, set
    Best-return-code to 5, "Downstream Mapping Mismatch". An
    Interface and Label Stack TLV SHOULD be included in the
    reply and filled in based on Interface-I and Stack-R. Go
    to step 7 (Send Reply Packet).

}

}

For each available downstream ECMP path {

    Retrieve output interface from the NHLFE entry.

    /* Note: this Return Code is set even if Label-stack-depth
    is one */

    If the output interface is not MPLS enabled {

        Set Best-return-code to Return Code 9, "Label switched
        but no MPLS forwarding at stack-depth" and set
        Best-rtn-subcode to Label-stack-depth and go to step 7
        (Send Reply Packet).

    }

    If a Downstream Detailed Mapping TLV is present {

        A Downstream Detailed Mapping TLV SHOULD be included in
        the echo reply (see Section 3.4) filled in with
        information about the current ECMP path.

    }

}

}
```

If no Downstream Detailed Mapping TLV is present, or the Downstream IP Address is set to the ALLROUTERS multicast address, go to step 7 (Send Reply Packet).

If the "Validate FEC Stack" flag is not set and the LSR is not configured to perform FEC checking by default, go to step 7 (Send Reply Packet).

```
/* Validate the Target FEC Stack in the received echo request.
```

```
First determine FEC-stack-depth from the Downstream Detailed Mapping TLV. This is done by walking through Stack-D (the Downstream labels) from the bottom, decrementing the number of labels for each non-Implicit Null label, while incrementing FEC-stack-depth for each label. If the Downstream Detailed Mapping TLV contains one or more Implicit Null labels, FEC-stack-depth may be greater than Label-stack-depth. To be consistent with the above stack-depths, the bottom is considered to be entry 1.
*/
```

```
Set FEC-stack-depth to 0. Set i to Label-stack-depth.
```

```
While (i > 0) do {
```

```
    ++FEC-stack-depth.
    if Stack-D [ FEC-stack-depth ] != 3 (Implicit Null)
        --i.
```

```
}
```

```
If the number of FECs in the FEC stack is greater than or equal to FEC-stack-depth {
Perform the FEC Checking procedure (see Section 4.4.1).
```

```
    If FEC-status is 2, set Best-return-code to 10 ("Mapping for this FEC is not the given label at stack-depth").
```

```
    If the Return Code is 1, set Best-return-code to FEC-return-code and Best-rtn-subcode to FEC-stack-depth.
```

```
}
```

```
Go to step 7 (Send Reply Packet).
```

```
}
```

5. Egress Processing:

```
/* These steps are performed by the LSR that identified itself as
the tail-end LSR for an LSP. */
```

```
If the received echo request contains no Downstream Detailed
Mapping TLV, or the Downstream IP Address is set to 127.0.0.1 or
0::1, go to step 6 (Egress FEC Validation).
```

```
Verify that the IP address, interface address, and label stack in
the Downstream Detailed Mapping TLV match Interface-I and Stack-R.
If not, set Best-return-code to 5, "Downstream Mapping Mismatch".
A Received Interface and Label Stack TLV SHOULD be created for the
echo response packet. Go to step 7 (Send Reply Packet).
```

6. Egress FEC Validation:

```
/* This is a loop for all entries in the Target FEC Stack starting
with FEC-stack-depth. */
```

```
Perform FEC checking by following the algorithm described in
Section 4.4.1 for Label-L and the FEC at FEC-stack-depth.
```

```
Set Best-return-code to FEC-code and Best-rtn-subcode to the value
in FEC-stack-depth.
```

```
If FEC-status (the result of the check) is 1,
go to step 7 (Send Reply Packet).
```

```
/* Iterate to the next FEC entry */
```

```
++FEC-stack-depth.
```

```
If FEC-stack-depth > the number of FECs in the FEC-stack,
go to step 7 (Send Reply Packet).
```

```
If FEC-status is 0 {
```

```
  ++Label-stack-depth.
```

```
  If Label-stack-depth > the number of labels in Stack-R,
  go to step 7 (Send Reply Packet).
```

```
  Label-L = extracted label from Stack-R at depth
```

```
  Label-stack-depth.
```

```
  Loop back to step 6 (Egress FEC Validation).
```

```
}
```

7. Send Reply Packet:

Send an MPLS echo reply with a Return Code of Best-return-code and a Return Subcode of Best-rtn-subcode. Include any TLVs created during the above process. The procedures for sending the echo reply are found in Section 4.5.

4.4.1. FEC Validation

```
/* This section describes validation of a FEC entry within the Target
FEC Stack and accepts a FEC, Label-L, and Interface-I.
```

```
If the outermost FEC of the Target FEC stack is the Nil FEC, then the
node MUST skip the Target FEC validation completely. This is to
support FEC hiding, in which the outer hidden FEC can be the Nil FEC.
Else, the algorithm performs the following steps. */
```

1. Two return values, FEC-status and FEC-return-code, are initialized to 0.
2. If the FEC is the Nil FEC {
 If Label-L is either Explicit_Null or Router_Alert, return.
 Else {
 Set FEC-return-code to 10 ("Mapping for this FEC is not the
 given label at stack-depth").
 Set FEC-status to 1
 Return.
 }
}
3. Check the FEC label mapping that describes how traffic received on the LSP is further switched or which application it is associated with. If no mapping exists, set FEC-return-code to Return 4, "Replying router has no mapping for the FEC at stack-depth". Set FEC-status to 1. Return.
4. If the label mapping for FEC is Implicit Null, set FEC-status to 2 and proceed to step 5. Otherwise, if the label mapping for FEC is Label-L, proceed to step 5. Otherwise, set FEC-return-code to 10 ("Mapping for this FEC is not the given label at stack-depth"), set FEC-status to 1, and return.

5. This is a protocol check. Check what protocol would be used to advertise the FEC. If it can be determined that no protocol associated with Interface-I would have advertised a FEC of that FEC-Type, set FEC-return-code to 12 ("Protocol not associated with interface at FEC stack-depth"). Set FEC-status to 1.
6. Return.

4.5. Sending an MPLS Echo Reply

An MPLS echo reply is a UDP packet. It MUST ONLY be sent in response to an MPLS echo request. The source IP address is a routable address of the replier; the source port is the well-known UDP port for LSP ping. The destination IP address and UDP port are copied from the source IP address and UDP port of the echo request. The IP TTL is set to 255. If the Reply Mode in the echo request is "Reply via an IPv4 UDP packet with Router Alert", then the IP header MUST contain the Router Alert IP Option of value 0x0 [RFC2113] for IPv4 or 69 [RFC7506] for IPv6. If the reply is sent over an LSP, the topmost label MUST in this case be the Router Alert label (1) (see [RFC3032]).

The format of the echo reply is the same as the echo request. The Sender's Handle, the Sequence Number, and TimeStamp Sent are copied from the echo request; the TimeStamp Received is set to the time of day that the echo request is received (note that this information is most useful if the time-of-day clocks on the requester and the replier are synchronized). The FEC Stack TLV from the echo request MAY be copied to the reply.

The replier MUST fill in the Return Code and Subcode, as determined in the previous section.

If the echo request contains a Pad TLV, the replier MUST interpret the first octet for instructions regarding how to reply.

If the replying router is the destination of the FEC, then Downstream Detailed Mapping TLVs SHOULD NOT be included in the echo reply.

If the echo request contains a Downstream Detailed Mapping TLV, and the replying router is not the destination of the FEC, the replier SHOULD compute its downstream routers and corresponding labels for the incoming label and add Downstream Detailed Mapping TLVs for each one to the echo reply it sends back. A replying node should follow the procedures defined in Section 4.5.1 if there is a FEC stack change due to tunneled LSP. If the FEC stack change is due to stitched LSP, it should follow the procedures defined in Section 4.5.2.

If the Downstream Detailed Mapping TLV contains Multipath Information requiring more processing than the receiving router is willing to perform, the responding router MAY choose to respond with only a subset of multipaths contained in the echo request Downstream Detailed Mapping. (Note: The originator of the echo request MAY send another echo request with the Multipath Information that was not included in the reply.)

Except in the case of Reply Mode 4, "Reply via application-level control channel", echo replies are always sent in the context of the IP/MPLS network.

4.5.1. Addition of a New Tunnel

A transit node knows when the FEC being traced is going to enter a tunnel at that node. Thus, it knows about the new outer FEC. All transit nodes that are the origination point of a new tunnel SHOULD add the FEC stack change sub-TLV (Section 3.4.1.3) to the Downstream Detailed Mapping TLV in the echo reply. The transit node SHOULD add one FEC stack change sub-TLV of operation type PUSH, per new tunnel being originated at the transit node.

A transit node that sends a Downstream FEC stack change sub-TLV in the echo reply SHOULD fill the address of the remote peer, which is the peer of the current LSP being traced. If the transit node does not know the address of the remote peer, it MUST set the address type to Unspecified.

The Label Stack sub-TLV MUST contain one additional label per FEC being PUSHed. The label MUST be encoded as defined in Section 3.4.1.2. The label value MUST be the value used to switch the data traffic. If the tunnel is a transparent pipe to the node, i.e., the data-plane trace will not expire in the middle of the new tunnel, then a FEC stack change sub-TLV SHOULD NOT be added, and the Label Stack sub-TLV SHOULD NOT contain a label corresponding to the hidden tunnel.

If the transit node wishes to hide the nature of the tunnel from the ingress of the echo request, then it MAY not want to send details about the new tunnel FEC to the ingress. In such a case, the transit node SHOULD use the Nil FEC. The echo reply would then contain a FEC stack change sub-TLV with operation type PUSH and a Nil FEC. The value of the label in the Nil FEC MUST be set to zero. The remote peer address type MUST be set to Unspecified. The transit node SHOULD add one FEC stack change sub-TLV of operation type PUSH, per new tunnel being originated at the transit node. The Label Stack sub-TLV MUST contain one additional label per FEC being PUSHed. The label value MUST be the value used to switch the data traffic.

4.5.2. Transition between Tunnels

A transit node stitching two LSPs SHOULD include two FEC stack change sub-TLVs. One with a pop operation for the old FEC (ingress) and one with the PUSH operation for the new FEC (egress). The replying node SHOULD set the Return Code to "Label switched with FEC change" to indicate change in the FEC being traced.

If the replying node wishes to perform FEC hiding, it SHOULD respond back with two FEC stack change sub-TLVs, one pop followed by one PUSH. The pop operation MAY either exclude the FEC TLV (by setting the FEC TLV length to 0) or set the FEC TLV to contain the LDP FEC. The PUSH operation SHOULD have the FEC TLV containing the Nil FEC. The Return Code SHOULD be set to "Label switched with FEC change".

If the replying node wishes to perform FEC hiding, it MAY choose to not send any FEC stack change sub-TLVs in the echo reply if the number of labels does not change for the downstream node and the FEC type also does not change (Nil FEC). In such case, the replying node MUST NOT set the Return Code to "Label switched with FEC change".

4.6. Receiving an MPLS Echo Reply

An LSR X should only receive an MPLS echo reply in response to an MPLS echo request that it sent. Thus, on receipt of an MPLS echo reply, X should parse the packet to ensure that it is well-formed, then attempt to match up the echo reply with an echo request that it had previously sent, using the destination UDP port and the Sender's Handle. If no match is found, then X jettisons the echo reply; otherwise, it checks the Sequence Number to see if it matches.

If the echo reply contains Downstream Detailed Mappings, and X wishes to traceroute further, it SHOULD copy the Downstream Detailed Mapping(s) into its next echo request(s) (with TTL incremented by one).

If one or more FEC stack change sub-TLVs are received in the MPLS echo reply, the ingress node SHOULD process them and perform some validation.

The FEC stack changes are associated with a downstream neighbor and along a particular path of the LSP. Consequently, the ingress will need to maintain a FEC stack per path being traced (in case of multipath). All changes to the FEC stack resulting from the processing of a FEC stack change sub-TLV(s) should be applied only for the path along a given downstream neighbor. The following algorithm should be followed for processing FEC stack change sub-TLVs.


```
push_seen = FALSE
fec_stack_depth = current-depth-of-fec-stack-being-traced
saved_fec_stack = current_fec_stack

while (sub-tlv = get_next_sub_tlv(downstream_detailed_map_tlv))

    if (sub-tlv == NULL) break

    if (sub-tlv.type == FEC-Stack-Change) {

        if (sub-tlv.operation == POP) {
            if (push_seen) {
                Drop the echo reply
                current_fec_stack = saved_fec_stack
                return
            }

            if (fec_stack_depth == 0) {
                Drop the echo reply
                current_fec_stack = saved_fec_stack
                return
            }

            Pop FEC from FEC stack being traced
            fec_stack_depth--;
        }

        if (sub-tlv.operation == PUSH) {
            push_seen = 1
            Push FEC on FEC stack being traced
            fec_stack_depth++;
        }
    }
}

if (fec_stack_depth == 0) {
    Drop the echo reply
    current_fec_stack = saved_fec_stack
    return
}
```

The next MPLS echo request along the same path should use the modified FEC stack obtained after processing the FEC stack change sub-TLVs. A non-Nil FEC guarantees that the next echo request along the same path will have the Downstream Detailed Mapping TLV validated for IP address, interface address, and label stack mismatches.

If the top of the FEC stack is a Nil FEC and the MPLS echo reply does not contain any FEC stack change sub-TLVs, then it does not necessarily mean that the LSP has not started traversing a different tunnel. It could be that the LSP associated with the Nil FEC terminated at a transit node, and at the same time, a new LSP started at the same transit node. The Nil FEC would now be associated with the new LSP (and the ingress has no way of knowing this). Thus, it is not possible to build an accurate hierarchical LSP topology if a traceroute contains Nil FECs.

A reply from a downstream node with Return Code 3, may not necessarily be for the FEC being traced. It could be for one of the new FECs that was added. On receipt of an IS_EGRESS reply, the LSP ingress should check if the depth of Target FEC sent to the node that just responded was the same as the depth of the FEC that was being traced. If it was not, then it should pop an entry from the Target FEC stack and resend the request with the same TTL (as previously sent). The process of popping a FEC is to be repeated until either the LSP ingress receives a non-IS_EGRESS reply or until all the additional FECs added to the FEC stack have already been popped. Using an IS_EGRESS reply, an ingress can build a map of the hierarchical LSP structure traversed by a given FEC.

When the MPLS echo reply Return Code is "Label switched with FEC change", the ingress node SHOULD manipulate the FEC stack as per the FEC stack change sub-TLVs contained in the Downstream Detailed Mapping TLV. A transit node can use this Return Code for stitched LSPs and for hierarchical LSPs. In case of ECMP or P2MP, there could be multiple paths and Downstream Detailed Mapping TLVs with different Return Codes (see Section 3.1, Note 2). The ingress node should build the topology based on the Return Code per ECMP path/P2MP branch.

4.7. Issue with VPN IPv4 and IPv6 Prefixes

Typically, an LSP ping for a VPN IPv4 prefix or VPN IPv6 prefix is sent with a label stack of depth greater than 1, with the innermost label having a TTL of 1. This is to terminate the ping at the egress PE, before it gets sent to the customer device. However, under certain circumstances, the label stack can shrink to a single label before the ping hits the egress PE; this will result in the ping terminating prematurely. One such scenario is a multi-AS Carrier's Carrier VPN.

To get around this problem, one approach is for the LSR that receives such a ping to realize that the ping terminated prematurely and to send back Return Code 13. In that case, the initiating LSR can retry

the ping after incrementing the TTL on the VPN label. In this fashion, the ingress LSR will sequentially try TTL values until it finds one that allows the VPN ping to reach the egress PE.

4.8. Non-compliant Routers

If the egress for the FEC Stack being pinged does not support LSP ping, then no reply will be sent, resulting in possible "false negatives". When in "traceroute" mode, if a transit LSR does not support LSP ping, then no reply will be forthcoming from that LSR for some TTL, say, n . The LSR originating the echo request SHOULD try sending the echo request with TTL= $n+1$, $n+2$, ..., $n+k$ to probe LSRs further down the path. In such a case, the echo request for TTL $> n$ SHOULD be sent with the Downstream Detailed Mapping TLV "Downstream IP Address" field set to the ALLROUTERS multicast address until a reply is received with a Downstream Detailed Mapping TLV. The label Stack TLV MAY be omitted from the Downstream Detailed Mapping TLV. Furthermore, the "Validate FEC Stack" flag SHOULD NOT be set until an echo reply packet with a Downstream Detailed Mapping TLV is received.

5. Security Considerations

Overall, the security needs for LSP ping are similar to those of ICMP ping.

There are at least three approaches to attacking LSRs using the mechanisms defined here. One is a Denial-of-Service (DoS) attack, by sending MPLS echo requests/replies to LSRs and thereby increasing their workload. The second is obfuscating the state of the MPLS data-plane liveness by spoofing, hijacking, replaying, or otherwise tampering with MPLS echo requests and replies. The third is an unauthorized source using an LSP ping to obtain information about the network.

To avoid potential DoS attacks, it is RECOMMENDED that implementations regulate the LSP ping traffic going to the control plane. A rate limiter SHOULD be applied to the well-known UDP port defined in Section 6.1.

Unsophisticated replay and spoofing attacks involving faking or replaying MPLS echo reply messages are unlikely to be effective. These replies would have to match the Sender's Handle and Sequence Number of an outstanding MPLS echo request message. A non-matching replay would be discarded as the sequence has moved on, thus a spoof has only a small window of opportunity. However, to provide a stronger defense, an implementation MAY also validate the TimeStamp Sent by requiring an exact match on this field.

To protect against unauthorized sources using MPLS echo request messages to obtain network information, it is RECOMMENDED that implementations provide a means of checking the source addresses of MPLS echo request messages against an access list before accepting the message.

It is not clear how to prevent hijacking (non-delivery) of echo requests or replies; however, if these messages are indeed hijacked, LSP ping will report that the data plane is not working as it should.

It does not seem vital (at this point) to secure the data carried in MPLS echo requests and replies, although knowledge of the state of the MPLS data plane may be considered confidential by some. Implementations SHOULD, however, provide a means of filtering the addresses to which echo reply messages may be sent.

The value part of the Pad TLV contains a variable number of octets. With the exception of the first octet, these contents, if any, are ignored on receipt, and can therefore serve as a clandestine channel.

When MPLS LSP ping is used within an administrative domain, a deployment can increase security by using border filtering of incoming LSP ping packets as well as outgoing LSP ping packets.

Although this document makes special use of 127/8 addresses, these are used only in conjunction with the UDP port 3503. Furthermore, these packets are only processed by routers. All other hosts MUST treat all packets with a destination address in the range 127/8 in accordance to RFC 1122. Any packet received by a router with a destination address in the range 127/8 without a destination UDP port of 3503 MUST be treated in accordance to RFC 1812. In particular, the default behavior is to treat packets destined to a 127/8 address as "martians".

If a network operator wants to prevent tracing inside a tunnel, one can use the Pipe Model [RFC3443], i.e., hide the outer MPLS tunnel by not propagating the MPLS TTL into the outer tunnel (at the start of the outer tunnel). By doing this, LSP traceroute packets will not expire in the outer tunnel, and the outer tunnel will not get traced.

If one doesn't wish to expose the details of the new outer LSP, then the Nil FEC can be used to hide those details. Using the Nil FEC ensures that the trace progresses without false negatives and all transit nodes (of the new outer tunnel) perform some minimal validations on the received MPLS echo requests.

6. IANA Considerations

6.1. TCP and UDP Port Number

The TCP and UDP port number 3503 has been allocated by IANA for LSP echo requests and replies.

6.2. MPLS LSP Ping Parameters

IANA maintains the "Multiprotocol Label Switching (MPLS) Label Switched Paths (LSPs) Ping Parameters" registry at [IANA-MPLS-LSP-PING].

The following subsections detail the name spaces managed by IANA. For some of these name spaces, the space is divided into assignment ranges; the following terms are used in describing the procedures by which IANA allocates values: "Standards Action" (as defined in [RFC5226]), "Specification Required", and "Vendor Private Use".

Values from "Specification Required" ranges MUST be registered with IANA. The request MUST be made via an RFC that describes the format and procedures for using the code point; the actual assignment is made during the IANA actions for the RFC.

Values from "Vendor Private" ranges MUST NOT be registered with IANA; however, the message MUST contain an enterprise code as registered with the IANA SMI Private Network Management Private Enterprise Numbers. For each name space that has a Vendor Private range, it must be specified where exactly the SMI Private Enterprise Number resides; see below for examples. In this way, several enterprises (vendors) can use the same code point without fear of collision.

6.2.1. Message Types, Reply Modes, Return Codes

IANA has created and will maintain registries for Message Types, Reply Modes, and Return Codes. Each of these can take values in the range 0-255. Assignments in the range 0-191 are via Standards Action; assignments in the range 192-251 are made via "Specification Required"; values in the range 252-255 are for Vendor Private Use and MUST NOT be allocated.

If any of these fields fall in the Vendor Private range, a top-level Vendor Enterprise Number TLV MUST be present in the message.

Message Types defined in this document are the following:

Value	Meaning
-----	-----
1	MPLS Echo Request
2	MPLS Echo Reply

Reply Modes defined in this document are the following:

Value	Meaning
-----	-----
1	Do not reply
2	Reply via an IPv4/IPv6 UDP packet
3	Reply via an IPv4/IPv6 UDP packet with Router Alert
4	Reply via application-level control channel

Return Codes defined in this document are listed in Section 3.1.

IANA has updated the reference for each these values to this document.

6.2.2. TLVs

IANA has created and maintains a registry for the Type field of top-level TLVs as well as for any associated sub-TLVs. Note that the meaning of a sub-TLV is scoped by the TLV. The number spaces for the sub-TLVs of various TLVs are independent.

The valid range for TLVs and sub-TLVs is 0-65535. Assignments in the ranges 0-16383 and 32768-49161 are made via Standards Action as defined in [RFC5226]; assignments in the ranges 16384-31743 and 49162-64511 are made via "Specification Required"; values in the ranges 31744-32767 and 64512-65535 are for Vendor Private Use and MUST NOT be allocated.

If a TLV or sub-TLV has a Type that falls in the range for Vendor Private Use, the Length MUST be at least 4, and the first four octets MUST be that vendor's SMI Private Enterprise Number, in network octet order. The rest of the Value field is private to the vendor.

TLVs and sub-TLVs defined in this document are the following:

Type	Sub-Type	Value Field
----	-----	-----
1		Target FEC Stack
	1	LDP IPv4 prefix
	2	LDP IPv6 prefix
	3	RSVP IPv4 LSP
	4	RSVP IPv6 LSP
	5	Unassigned
	6	VPN IPv4 prefix
	7	VPN IPv6 prefix
	8	L2 VPN endpoint
	9	"FEC 128" Pseudowire - IPv4 (Deprecated)
	10	"FEC 128" Pseudowire - IPv6
	11	"FEC 129" Pseudowire - IPv4
	12	BGP labeled IPv4 prefix
	13	BGP labeled IPv6 prefix
	14	Generic IPv4 prefix
	15	Generic IPv6 prefix
	16	Nil FEC
	24	"FEC 128" Pseudowire - IPv6
	25	"FEC 129" Pseudowire - IPv6
2		Downstream Mapping (Deprecated)
3		Pad
4		Unassigned
5		Vendor Enterprise Number
6		Unassigned
7		Interface and Label Stack
8		Unassigned
9		Errored TLVs
	Any value	The TLV not understood
10		Reply TOS Byte
20		Downstream Detailed Mapping

IANA has updated the reference for each of these values to this document.

6.2.3. Global Flags

IANA has created a "Global Flags" subregistry of the "Multiprotocol Label Switching (MPLS) Label Switched Paths (LSPs) Ping Parameters" registry.

This registry tracks the assignment of 16 flags in the Global Flags field of the MPLS LSP ping echo request message. The flags are numbered from 0 (most significant bit, transmitted first) to 15.

New entries are assigned by Standards Action.

Initial entries in the registry are as follows:

Bit number	Name	Reference
15	V Flag	[RFC8029]
14	T Flag	[RFC6425]
13	R Flag	[RFC6426]
12-0	Unassigned	[RFC8029]

6.2.4. Downstream Detailed Mapping Address Type

This document extends RFC 4379 by defining a new address type for use with the Downstream Mapping and Downstream Detailed Mapping TLVs. IANA has established a registry to assign address types for use with the Downstream Mapping and Downstream Detailed Mapping TLVs, which initially allocates the following assignments:

Type #	Address Type	K Octets	Reference
1	IPv4 Numbered	16	[RFC8029]
2	IPv4 Unnumbered	16	[RFC8029]
3	IPv6 Numbered	40	[RFC8029]
4	IPv6 Unnumbered	28	[RFC8029]
5	Non IP	12	[RFC6426]

Downstream Detailed Mapping Address Type Registry

Because the field in this case is an 8-bit field, the allocation policy for this registry is "Standards Action".

6.2.5. DS Flags

This document defines the Downstream Mapping (DSMAP) TLV and the Downstream Detailed Mapping (DDMAP) TLV, which have Type 2 and Type 20, respectively, assigned from the "TLVs" subregistry of the "Multiprotocol Label Switching (MPLS) Label Switched Paths (LSPs) Ping Parameters" registry.

DSMAP has been deprecated by DDMAP, but both TLVs share a field: DS Flags.

IANA has created and now maintains a registry entitled "DS Flags".

The registration policy for this registry is Standards Action [RFC5226].

IANA has made the following assignments:

Bit Number	Name	Reference
7	N: Treat as a Non-IP Packet	[RFC8029]
6	I: Interface and Label Stack Object Request	[RFC8029]
5	E: ELI/EL push indicator	[RFC8012]
4	L: Label-based load balance indicator	[RFC8012]
3-0	Unassigned	

6.2.6. Multipath Types

IANA has created and now maintains a registry entitled "Multipath Types".

The registration policy [RFC5226] for this registry is Standards Action.

IANA has made the following assignments:

Value	Meaning	Reference
0	no multipath	[RFC8029]
1	Unassigned	
2	IP address	[RFC8029]
3	Unassigned	
4	IP address range	[RFC8029]
5-7	Unassigned	
8	Bit-masked IP address set	[RFC8029]
9	Bit-masked label set	[RFC8029]
10	IP and label set	[RFC8012]
11-250	Unassigned	
251-254	Reserved for Experimental Use	[RFC8029]
255	Reserved	[RFC8029]

6.2.7. Pad Type

IANA has created and now maintains a registry entitled "Pad Types".

The registration policy [RFC5226] for this registry is Standards Action.

IANA has made the following initial assignments:

Registry Name: Pad Types

Value	Meaning	Reference
0	Reserved	[RFC8029]
1	Drop Pad TLV from reply	[RFC8029]
2	Copy Pad TLV to reply	[RFC8029]
3-250	Unassigned	
251-254	Experimental Use	[RFC8029]
255	Reserved	[RFC8029]

6.2.8. Interface and Label Stack Address Type

IANA has created and now maintains a registry entitled "Interface and Label Stack Address Types".

The registration policy [RFC5226] for this registry is Standards Action.

IANA has made the following initial assignments:

Registry Name: Interface and Label Stack Address Types

Value	Meaning	Reference
0	Reserved	[RFC8029]
1	IPv4 Numbered	[RFC8029]
2	IPv4 Unnumbered	[RFC8029]
3	IPv6 Numbered	[RFC8029]
4	IPv6 Unnumbered	[RFC8029]
5-250	Unassigned	
251-254	Experimental Use	[RFC8029]
255	Reserved	[RFC8029]

6.3. IPv4 Special-Purpose Address Registry

IANA has updated the reference in Note 1 of the "IANA IPv4 Special-Purpose Address Registry" [IANA-SPECIAL-IPv4] to point to this document.

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7.2. Informative References

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Appendix A. Deprecated TLVs and Sub-TLVs (Non-normative)

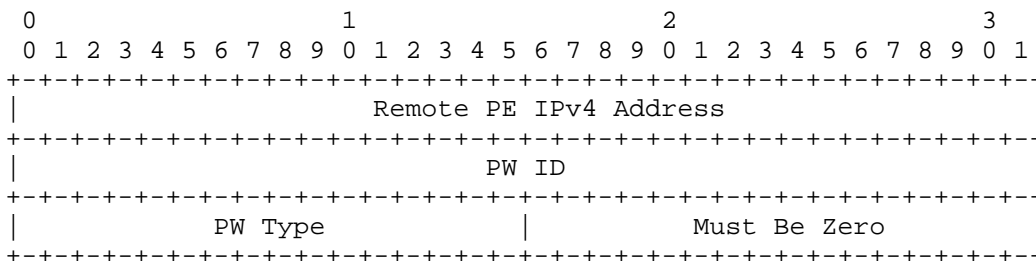
This appendix describes deprecated elements, which are non-normative for an implementation. They are included in this document for historical and informational purposes.

A.1. Target FEC Stack

A.1.1. FEC 128 Pseudowire - IPv4 (Deprecated)

FEC 128 (0x80) is defined in [RFC4447], as are the terms PW ID (Pseudowire ID) and PW Type (Pseudowire Type). A PW ID is a non-zero 32-bit connection ID. The PW Type is a 15-bit number indicating the encapsulation type. It is carried right justified in the field below termed encapsulation type with the high-order bit set to zero. Both of these fields are treated in this protocol as opaque values.

When a FEC 128 is encoded in a label stack, the following format is used. The Value field consists of the Remote PE IPv4 Address (the destination address of the targeted LDP session), the PW ID, and the encapsulation type as follows:



This FEC is deprecated and is retained only for backward compatibility. Implementations of LSP ping SHOULD accept and process this TLV, but SHOULD send LSP ping echo requests with the new TLV (see Section 3.2.9), unless explicitly configured to use the old TLV.

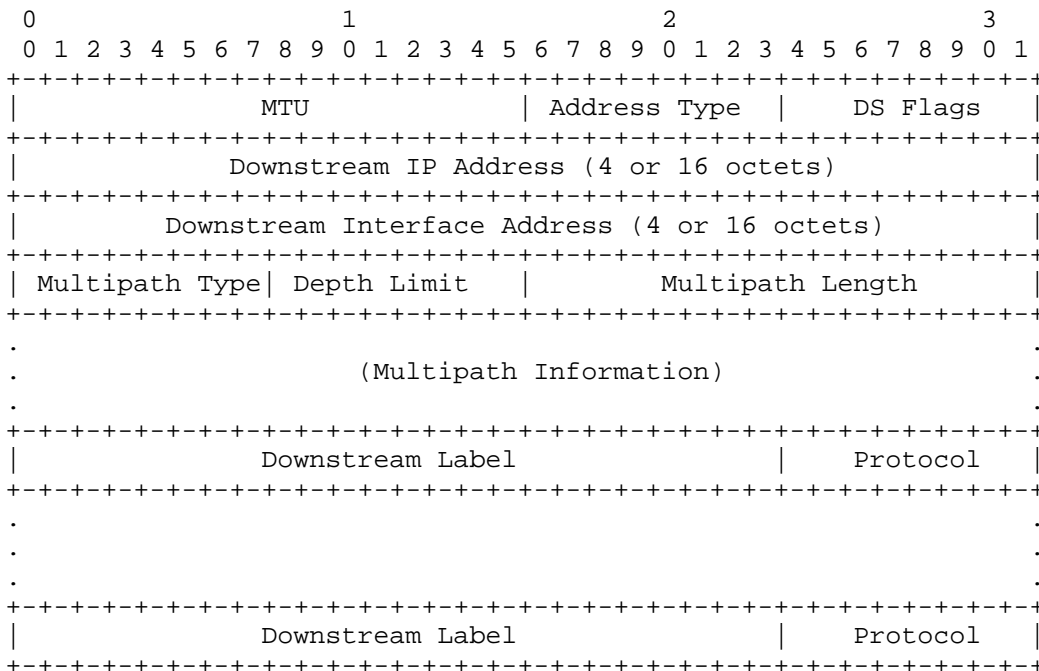
An LSR receiving this TLV SHOULD use the source IP address of the LSP echo request to infer the sender's PE address.

A.2. Downstream Mapping (Deprecated)

The Downstream Mapping object is a TLV that MAY be included in an echo request message. Only one Downstream Mapping object may appear in an echo request. The presence of a Downstream Mapping object is a request that Downstream Mapping objects be included in the echo reply. If the replying router is the destination of the FEC, then a Downstream Mapping TLV SHOULD NOT be included in the echo reply.

Otherwise, the replying router SHOULD include a Downstream Mapping object for each interface over which this FEC could be forwarded. For a more precise definition of the notion of "downstream", see Section 3.4.2, "Downstream Router and Interface".

The Length is K + M + 4*N octets, where M is the Multipath Length, and N is the number of downstream labels. Values for K are found in the description of Address Type below. The Value field of a Downstream Mapping has the following format:



Maximum Transmission Unit (MTU)

The MTU is the size in octets of the largest MPLS frame (including label stack) that fits on the interface to the downstream LSR.

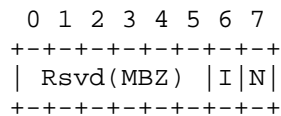
Address Type

The Address Type indicates if the interface is numbered or unnumbered. It also determines the length of the Downstream IP Address and Downstream Interface fields. The resulting total for the initial part of the TLV is listed in the table below as "K Octets". The Address Type is set to one of the following values:

Type #	Address Type	K Octets
-----	-----	-----
1	IPv4 Numbered	16
2	IPv4 Unnumbered	16
3	IPv6 Numbered	40
4	IPv6 Unnumbered	28
5	Non IP	12

DS Flags

The DS Flags field is a bit vector with the following format:



Two flags are defined currently, I and N. The remaining flags MUST be set to zero when sending and ignored on receipt.

Flag Name and Meaning

I Interface and Label Stack Object Request

When this flag is set, it indicates that the replying router SHOULD include an Interface and Label Stack Object in the echo reply message.

N Treat as a Non-IP Packet

Echo request messages will be used to diagnose non-IP flows. However, these messages are carried in IP packets. For a router that alters its ECMP algorithm based on the FEC or deep packet examination, this flag requests that the router treat this as it would if the determination of an IP payload had failed.

Downstream IP Address and Downstream Interface Address

IPv4 addresses and interface indices are encoded in 4 octets; IPv6 addresses are encoded in 16 octets.

If the interface to the downstream LSR is numbered, then the Address Type MUST be set to IPv4 or IPv6, the Downstream IP Address MUST be set to either the downstream LSR's Router ID or the interface address of the downstream LSR, and the Downstream Interface Address MUST be set to the downstream LSR's interface address.

If the interface to the downstream LSR is unnumbered, the Address Type MUST be IPv4 Unnumbered or IPv6 Unnumbered, the Downstream IP Address MUST be the downstream LSR's Router ID, and the Downstream Interface Address MUST be set to the index assigned by the upstream LSR to the interface.

If an LSR does not know the IP address of its neighbor, then it MUST set the Address Type to either IPv4 Unnumbered or IPv6 Unnumbered. For IPv4, it must set the Downstream IP Address to 127.0.0.1; for IPv6, the address is set to 0::1. In both cases, the interface index MUST be set to 0. If an LSR receives an Echo Request packet with either of these addresses in the Downstream IP Address field, this indicates that it MUST bypass interface verification but continue with label validation.

If the originator of an echo request packet wishes to obtain Downstream Mapping information but does not know the expected label stack, then it SHOULD set the Address Type to either IPv4 Unnumbered or IPv6 Unnumbered. For IPv4, it MUST set the Downstream IP Address to 224.0.0.2; for IPv6, the address MUST be set to FF02::2. In both cases, the interface index MUST be set to 0. If an LSR receives an echo request packet with the all-routers multicast address, then this indicates that it MUST bypass both interface and label stack validation, but return Downstream Mapping TLVs using the information provided.

Multipath Type

The following Multipath Types are defined:

Key	Type	Multipath Information
----	-----	-----
0	no multipath	Empty (Multipath Length = 0)
2	IP address	IP addresses
4	IP address range	low/high address pairs
8	Bit-masked IP address set	IP address prefix and bit mask
9	Bit-masked label set	Label prefix and bit mask

Type 0 indicates that all packets will be forwarded out this one interface.

Types 2, 4, 8, and 9 specify that the supplied Multipath Information will serve to exercise this path.

Depth Limit

The Depth Limit is applicable only to a label stack and is the maximum number of labels considered in the hash; this SHOULD be set to zero if unspecified or unlimited.

Multipath Length

The length in octets of the Multipath Information.

Multipath Information

Address or label values encoded according to the Multipath Type. See Section 3.4.1.1.1 for encoding details.

Downstream Label(s)

The set of labels in the label stack as it would have appeared if this router were forwarding the packet through this interface. Any Implicit Null labels are explicitly included. Labels are treated as numbers, i.e., they are right justified in the field.

A downstream label is 24 bits, in the same format as an MPLS label minus the TTL field, i.e., the MSBit of the label is bit 0, the LSBit is bit 19, the TC bits are bits 20-22, and bit 23 is the S bit. The replying router SHOULD fill in the TC and S bits; the LSR receiving the echo reply MAY choose to ignore these bits.

Protocol

The protocol is taken from the following table:

Protocol #	Signaling Protocol
-----	-----
0	Unknown
1	Static
2	BGP
3	LDP
4	RSVP-TE

Acknowledgements

The original acknowledgements from RFC 4379 state the following:

This document is the outcome of many discussions among many people, including Manoj Leelanivas, Paul Traina, Yakov Rekhter, Der-Hwa Gan, Brook Bailey, Eric Rosen, Ina Minei, Shivani Aggarwal, and Vanson Lim.

The description of the Multipath Information sub-field of the Downstream Mapping TLV was adapted from text suggested by Curtis Villamizar.

We would like to thank Loa Andersson for motivating the advancement of this specification.

We also would like to thank Alexander Vainshtein, Yimin Shen, Curtis Villamizar, David Allan, Vincent Roca, Mirja Kuhlewind, and Elwyn Davies for their review and useful comments.

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A mechanism used to detect data-plane failures in MPLS LSPs was originally published as RFC 4379 in February 2006. It was produced by the MPLS Working Group of the IETF and was jointly authored by Kireeti Kompella and George Swallow.

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